

# A First Investigation of Sturmian Trees

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# Outline

- 1 Sturmian words
- 2 Sturmian trees
- 3 Slow automata
- 4 Rank and degree
- 5 Results



# Factors of an infinite word

A factor  $w$  of a word  $x$  is a finite word that occurs in  $x$ , that is, there are words  $u$  and  $y$  such that  $x = uwy$ .

## Example (Fibonacci word)

$x =$  *a b a a b a b a a b a a b a b a a b a b a a b ...*

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Factors	<i>a</i>	<i>aa</i>	<i>aab</i>	<i>aaba</i>	<i>aabaa</i>
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Length: $n$	1	2	3	4	5
Factors	$a$	$aa$	$aab$	$aaba$	$aabaa$
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# Sturmian words

## Proposition (Hedlund & Morse)

*An infinite word  $x$  ultimately periodic iff there is an integer  $n$  such that  $x$  has at most  $n$  distinct factors of length  $n$ .*

An infinite word  $x$  is **Sturmian** if the number of its factors of length  $n$  is  $n + 1$  for each  $n$ .

Sturmian words are non ultimately periodic words with the smallest complexity.

Example (Fibonacci word:  $f_{n+2} = f_{n+1}f_n$ )

$$f_0 = a$$

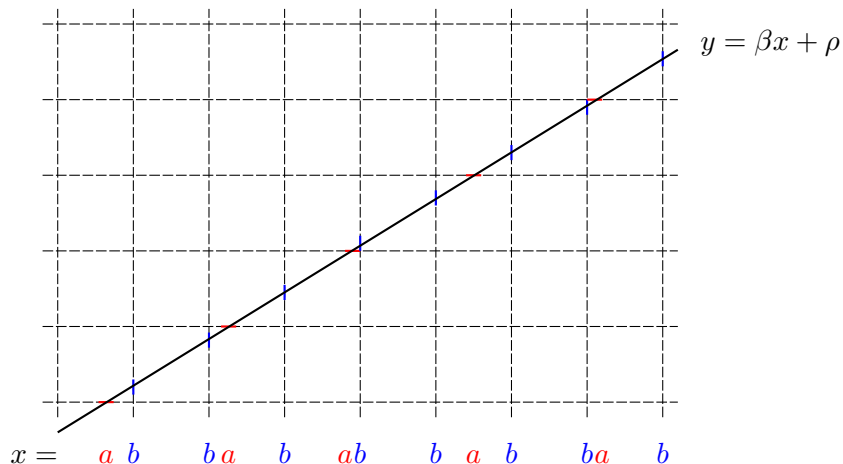
$$f_1 = ab$$

$$f_2 = aba$$

$$f_3 = abaab$$

$$f_\omega = abaababa \dots$$

# Characterization: cutting sequences

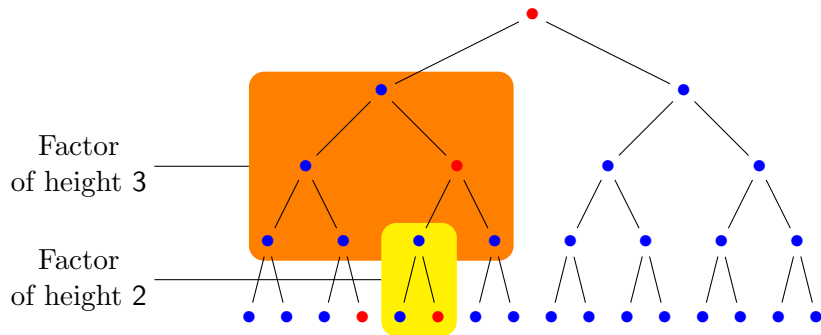


## Theorem

*A infinite word is Sturmian iff it is the cutting sequence of a straight line  $y = \beta x + \rho$  with an irrational slope  $\beta$ .*

# Factor of a tree

A **factor** of height  $h$  of a tree  $t$  is a subtree of height  $h$  that occurs in  $t$ .



## Sturmian tree

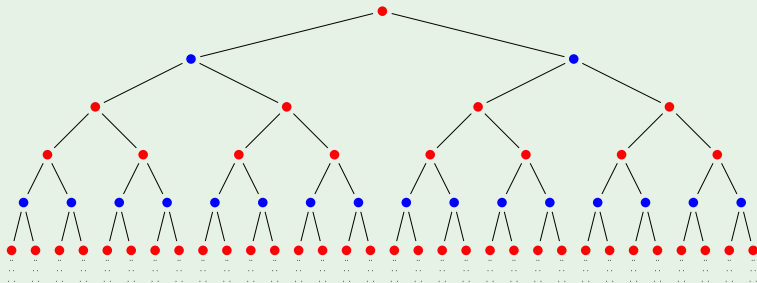
### Proposition (Carpi et al)

*A complete tree  $t$  is rational if there is some integer  $h$  such that  $t$  has at most  $h$  distinct factors of height  $h$ .*

A tree is **Sturmian** if the number of its factors of height  $h$  is  $h + 1$  for each  $h$ .

### Example (Easy one: uniform tree)

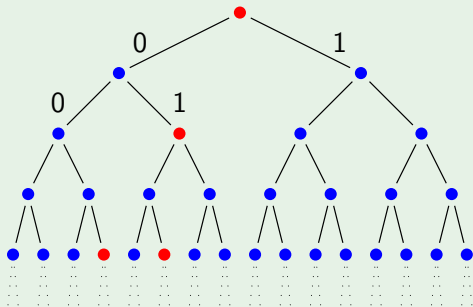
An Sturmian word  $x = \textcolor{red}{a}\textcolor{blue}{b}\textcolor{red}{a}\textcolor{blue}{b}\textcolor{red}{a}\dots$  is repeated on each branch.



Example (Unexpected one: Dyck tree)

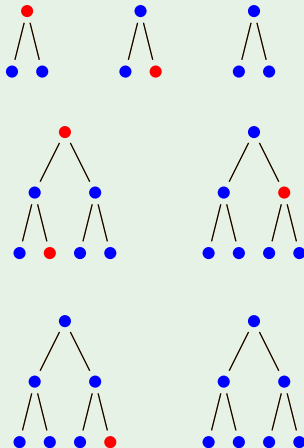
A node is  $\bullet$  if it is a Dyck word over the alphabet  $\{0, 1\}$ .

## The Dyck tree



$$D_2^* = \{\varepsilon, 01, 0101, 0011, \dots\}$$

Its factors

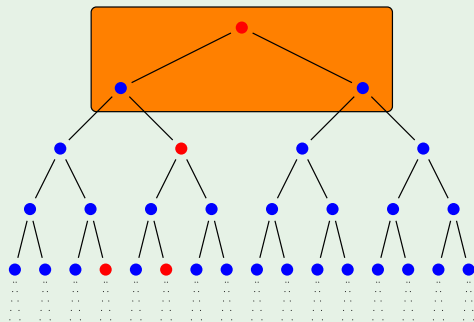




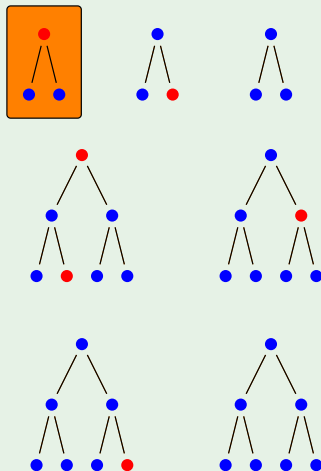
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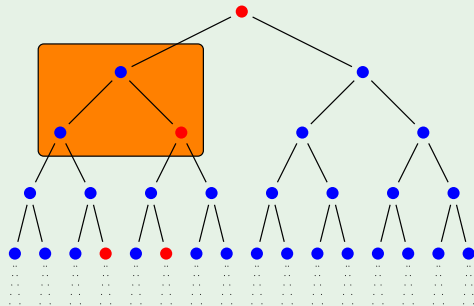
Its factors



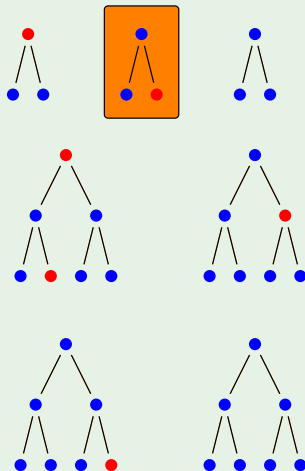
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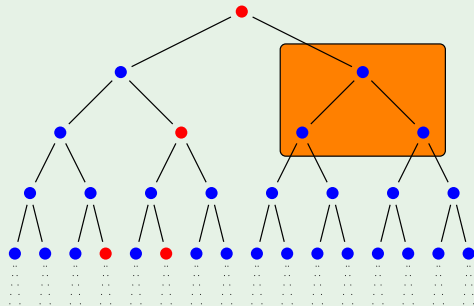
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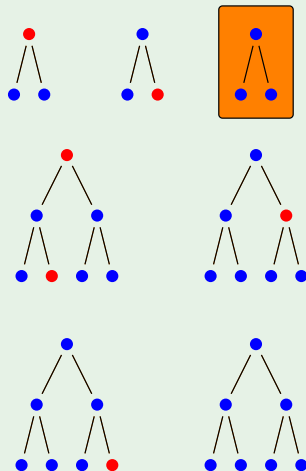
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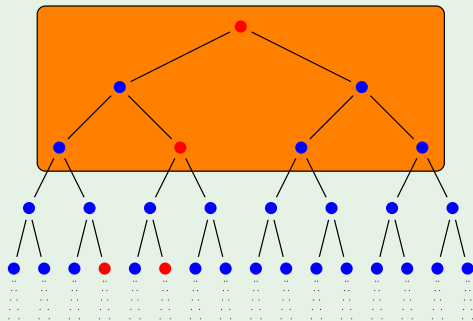
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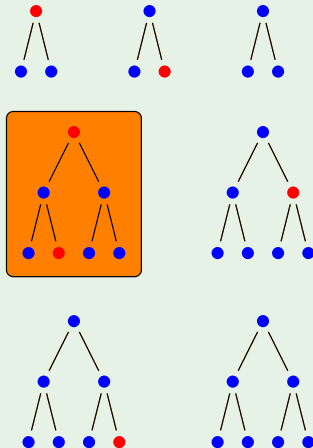
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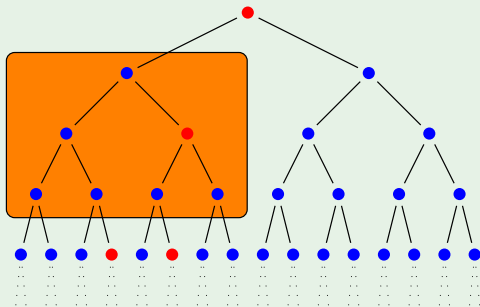
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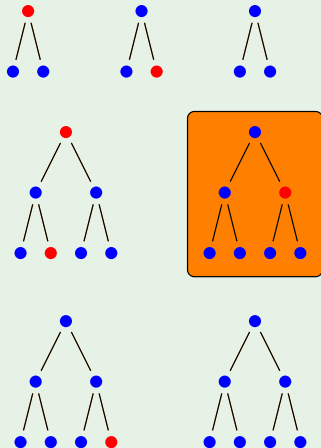
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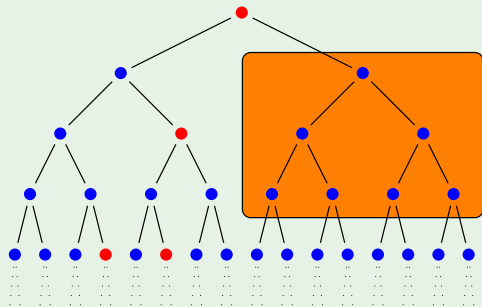
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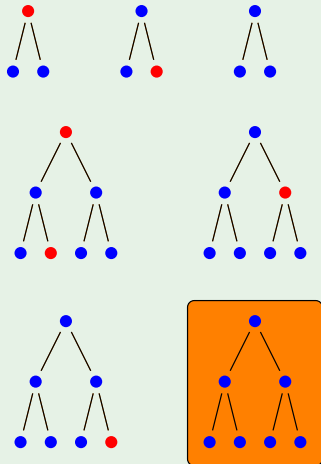
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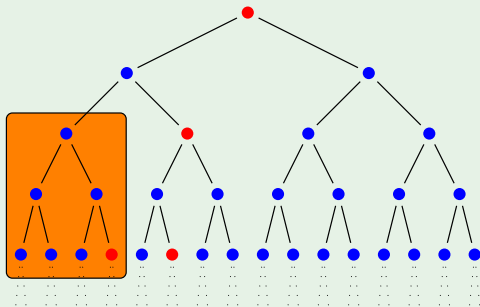
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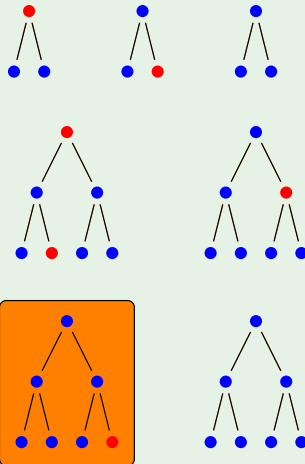
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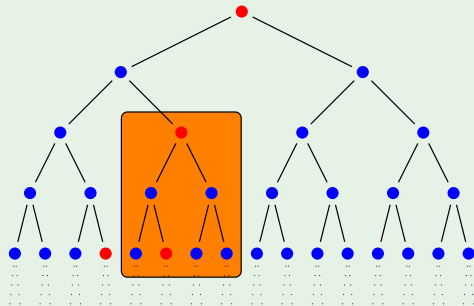
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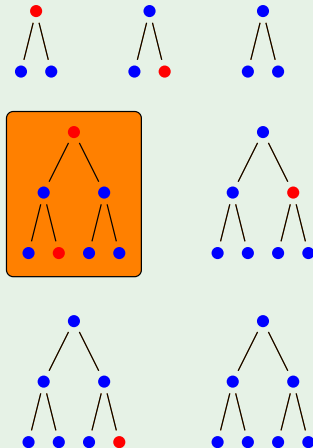
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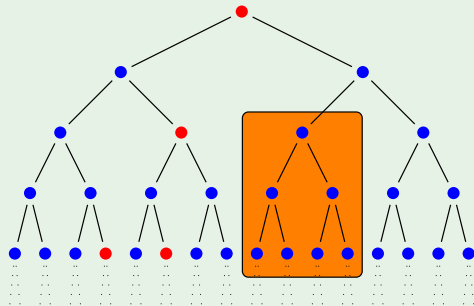




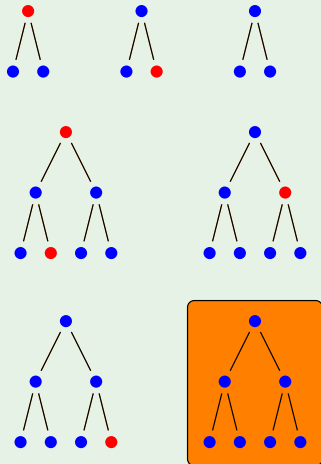
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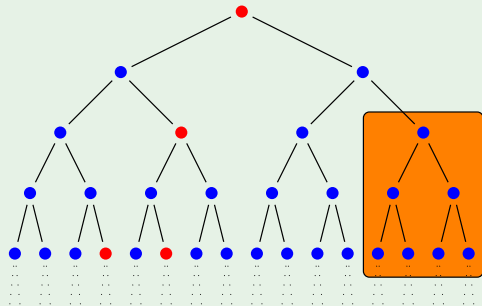
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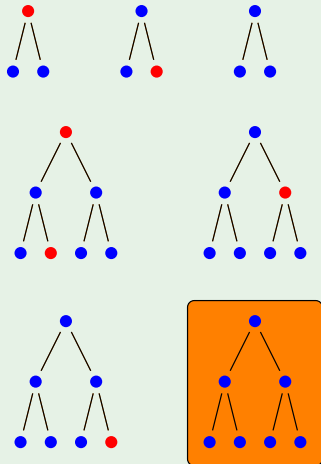
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# Slow automata

Let  $\mathcal{A} = (Q, A, \cdot, q_0, F)$  be a (infinite) deterministic automaton. Define the **Moore equivalence**  $\sim_n$  by induction.

$$\begin{aligned} q \sim_1 q' &\iff (q \in F \iff q' \in F) \\ q \sim_{n+1} q' &\iff (q \sim_n q') \text{ and } (\forall a \in A \quad q \cdot a \sim_n q' \cdot a) \end{aligned}$$

The relation  $q \sim_n q'$  does not hold if there is a word  $w$  of length  $n$  such that  $q \cdot w \in F$  and  $q' \cdot w \notin F$  (or  $q \cdot w \notin F$  and  $q' \cdot w \in F$ )

An infinite automaton is **slow** iff each equivalence  $\sim_n$  has  $n + 1$  classes.

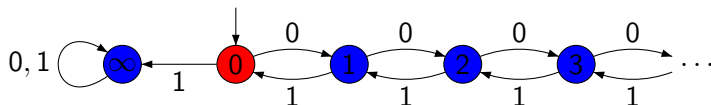
## Proposition

*A tree  $t$  is Sturmian iff the minimal automaton of  $t^{-1}(a)$  is slow.*



# Application to the Dyck tree

The minimal automaton of the Dyck language is the following.



The Moore equivalences of this automaton

$$\sim_1: \text{0} \mid \text{1, 2, 3, 4, } \dots \infty$$

$$\sim_2: \text{0} \mid \text{1} \mid \text{234 } \dots \infty$$

$$\sim_3: \text{0} \mid \text{1} \mid \text{2} \mid \text{3, 4, } \dots \infty$$

$$\sim_4: \text{0} \mid \text{1} \mid \text{2} \mid \text{3} \mid \text{4, } \dots \infty$$



# Rank and degree

- A node is called **irrational** if the infinite subtree rooted in this node is not rational.
- The **rank** is the number of distinct rational subtrees.
- The **degree** is the number of branches of irrational nodes.

## Examples

- The uniform tree has rank 0 and degree  $\infty$ .
- The Dyck tree has rank 1 and degree  $\infty$ .



# Results

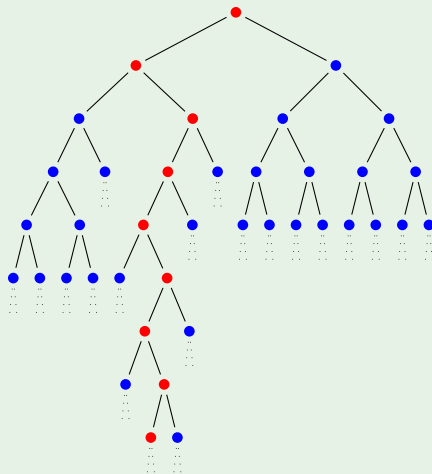
degree	rank	
	finite	infinite
1	characterized	example later
$\geq 2$ , finite	proved to be empty	example in full paper
infinite	example of Dyck tree	example in full paper



## Characterization

### Example (Indicator tree)

Take any Sturmian word (e.g. 01001010 $\dots$ ) and distinguish the **branch** labeled by this word.



Rank 1

Degree 1

# Yet another example

## Example (Rank $\infty$ and degree 1)

