Combinatorics of braids

Vincent Jugé

Université Paris Diderot (IRIF)

05/02/2016









Some questions of interest...





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What are braids?





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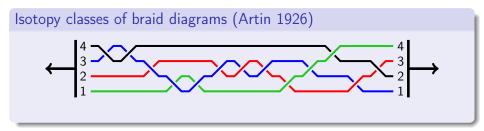
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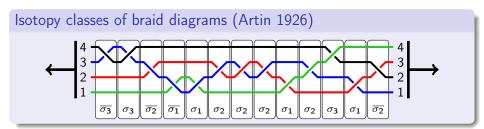
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- What is a complicated braid? Define notions of complexity.
- Mow do complicated braids typically behave?



Some questions of interest...

- What are braids? Mathematical objects interacting with each other.
- What is a complicated braid? Define notions of complexity.
- How do complicated braids typically behave? Choose a dynamic framework/probability measure.







Finitely generated monoid

$$\mathcal{B}_n = \left\langle \right.$$

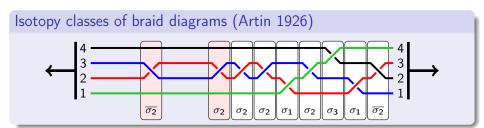


Isotopy classes of braid diagrams (Artin 1926)



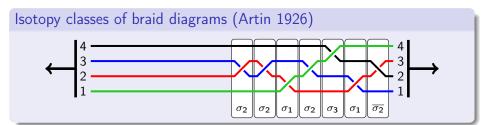
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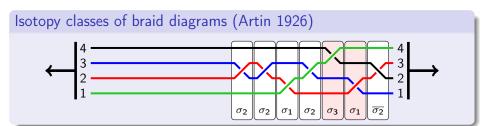
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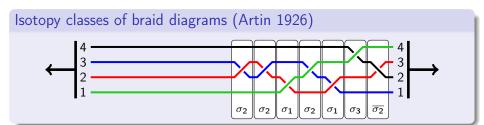


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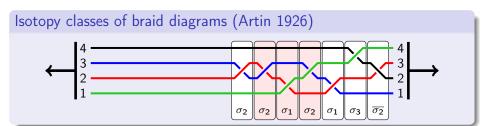




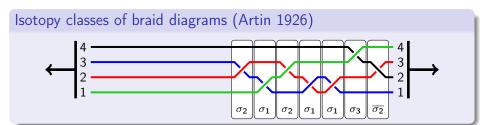
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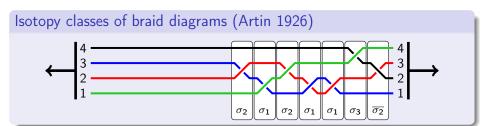
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Finitely presented group

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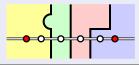
Isotopy classes of laminations of the punctured plane (Birman 1975)

Trivial lamination



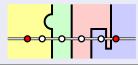
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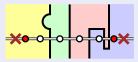
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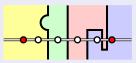
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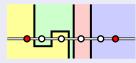


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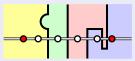


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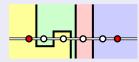


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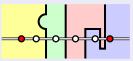
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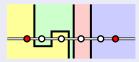


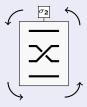
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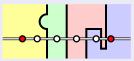
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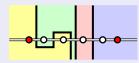


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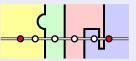
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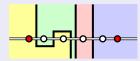


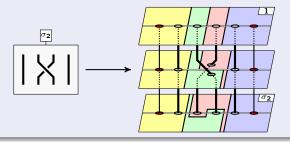
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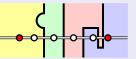
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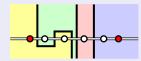


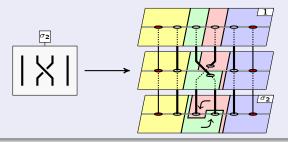
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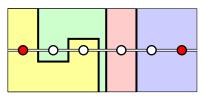
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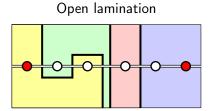


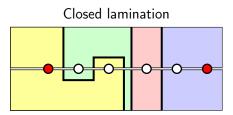
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Open lamination

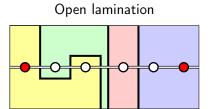


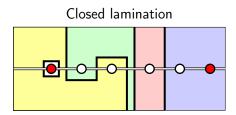
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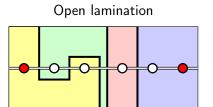


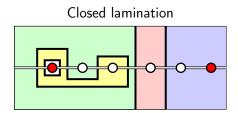
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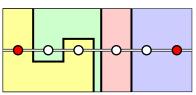
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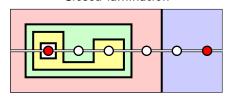


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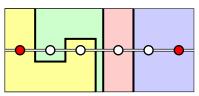


Closed lamination

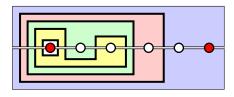


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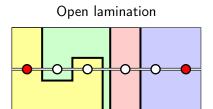
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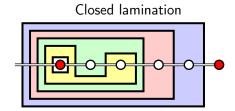


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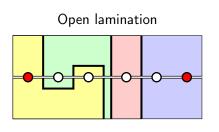


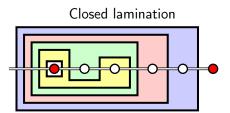
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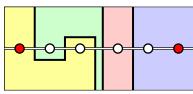




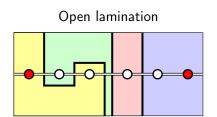
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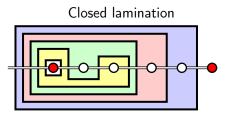


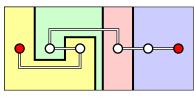




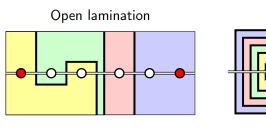
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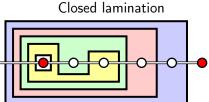


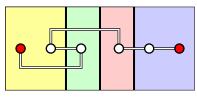




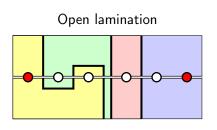
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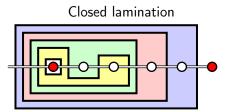


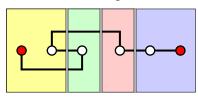




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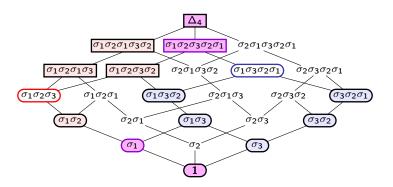


Garside normal form (Garside 1969, Adian 1984)

1 The monoid of **positive** braids $\mathcal{B}_n^+ = \langle \sigma_1, \dots, \sigma_{n-1} \rangle^+$ is a **lattice** for the divisibility ordering \leq . $(\alpha \leq \beta \Leftrightarrow \exists \gamma \in \mathcal{B}_n^+, \alpha \gamma = \beta)$

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- **2** There exists a Garside element $\Delta_n = \bigvee {\{\sigma_1, \dots, \sigma_{n-1}\}}$.
- **3** The Garside normal form of a positive braid $\alpha \in \mathcal{B}_n^+$ is the smallest word $Gar(\alpha) = a_1 \cdot a_2 \cdot \ldots \cdot a_k$ such that:

$$\alpha = a_1 a_2 \dots a_k; \qquad \qquad a_i = \Delta_n \wedge ((a_1 \dots a_{i-1})^{-1} \alpha).$$

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The Garside normal form:

- can be extended to the group \mathcal{B}_n ;
- is automatic: for all $i \in \{1, \ldots, n-1\}$, the languages $\{(\mathbf{Gar}(\alpha), \mathbf{Gar}(\alpha\sigma_i)) : \alpha \in \mathcal{B}_n\}$; $\{(\mathbf{Gar}(\alpha), \mathbf{Gar}(\sigma_i\alpha)) : \alpha \in \mathcal{B}_n\}$ are regular;
- solves the equality problem: $\alpha = \beta$ iff $Gar(\alpha) = Gar(\beta)$.

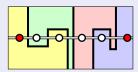
Tight laminations/curve diagrams

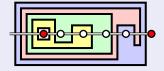
A lamination/curve diagram is **tight** if it minimises crossings + or + ·

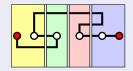
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Non-tight laminations/curve diagrams



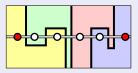


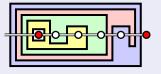


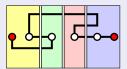
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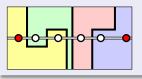
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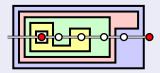


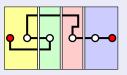




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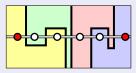


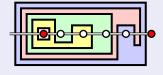


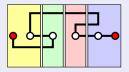
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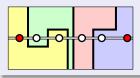
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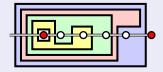


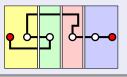




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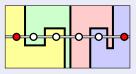


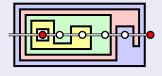
 Two tight laminations/curve diagrams represent the same braid iff they are visibly isotopic to each other. (≈ normal form)

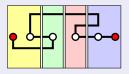
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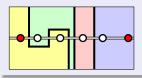
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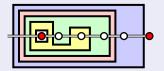


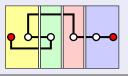




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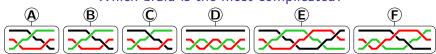




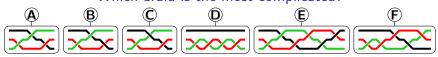


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Which braid is the most complicated?



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Several approaches to braid complexity

• "Naive" Artin length: $E > F > D > A \approx B \approx C$;

•
$$\mathbf{A} = \sigma_1 \sigma_2 \sigma_1$$
:

$$|A| = 3;$$

•
$$D = \sigma_1^4$$
:

$$|\mathbf{D}| = 4;$$

•
$$B = \sigma_1 \sigma_2 \overline{\sigma_1}$$
:

$$|B| = 3;$$

•
$$\mathbf{E} = (\sigma_1 \sigma_2)^3$$
:

$$|E| = 6;$$

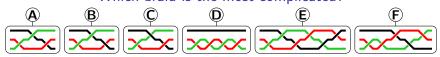
•
$$\mathbf{C} = \sigma_1 \overline{\sigma_2} \sigma_1$$
:

$$|\mathbf{C}| = 3$$
:

•
$$\mathbf{F} = \sigma_1^2 \sigma_2 \overline{\sigma_1 \sigma_2}$$
:

$$|F| = 5.$$

Which braid is the most complicated?



- "Naive" Artin length: $E > F > D > A \approx B \approx C$;
- "Real" Artin length: $E > D > A \approx B \approx C = F$;

•
$$A = \sigma_1 \sigma_2 \sigma_1$$
:

$$|A| = 3;$$

•
$$D = \sigma_1^4$$
:

$$|\mathbf{D}| = 4;$$

•
$$B = \sigma_1 \sigma_2 \overline{\sigma_1}$$
:

$$|B| = 3;$$

•
$$E = (\sigma_1 \sigma_2)^3$$
:

$$|E| = 6;$$

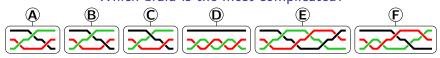
•
$$\mathbf{C} = \sigma_1 \overline{\sigma_2} \sigma_1$$
:

$$|\mathbf{C}| = 3;$$

•
$$\mathbf{F} = \sigma_1 \overline{\sigma_2} \sigma_1$$
:

$$|F| = 3.$$

Which braid is the most complicated?



- "Naive" Artin length: $E > F > D > A \approx B \approx C$;
- "Real" Artin length: $E > D > A \approx B \approx C = F$;
- Symmetric Garside length: $D > C = F > B \approx E > A$;

•
$$A = \sigma_1 \sigma_2 \sigma_1$$
: $|A| = 1$;

•
$$\mathbf{D} = \sigma_1 \cdot \sigma_1 \cdot \sigma_1 \cdot \sigma_1$$
: $|\mathbf{D}| = 4$;

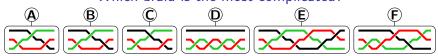
•
$$\mathbf{B} = \overline{\sigma_1} \cdot \sigma_1 \sigma_2$$
: $|\mathbf{B}| = 2$;

•
$$\mathbf{E} = \Delta_3 \cdot \Delta_3$$
: $|\mathbf{E}| = 2$;

•
$$\mathbf{C} = \overline{\sigma_2 \sigma_1} \cdot \sigma_2 \sigma_1 \cdot \sigma_1$$
: $|\mathbf{C}| = 3$;

•
$$F = C$$
: $|F| = 3$.

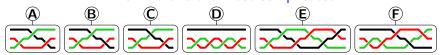
Which braid is the most complicated?



- "Naive" Artin length: $E > F > D > A \approx B \approx C$;
- "Real" Artin length: $E > D > A \approx B \approx C = F$;
- Symmetric Garside length: $D > C = F > B \approx E > A$;
- Open laminated complexity: $C = F > D \approx E > B > A$;

$$|A| = 6$$
 $|B| = 8$ $|C| = |F| = 14$ $|D| = 10$ $|E| = 10$

Which braid is the most complicated?



- "Naive" Artin length: $E > F > D > A \approx B \approx C$;
- $\bullet \text{ "Real" Artin length:} \qquad \qquad \mathsf{E} > \mathsf{D} > \mathsf{A} \approx \mathsf{B} \approx \mathsf{C} = \mathsf{F};$
- Symmetric Garside length: $D > C = F > B \approx E > A$;
- Open laminated complexity: $C = F > D \approx E > B > A$;
- Diagrammatic complexity: $C = F > D \approx E > B > A$.

$$|\mathbf{A}| = 6$$

$$|B| = 8$$

$$|C| = |F| = 14$$

$$|D| = 10$$

$$|{\bf E}| = 10$$











How fast can you compute the complexity of a braid $\alpha \in \mathcal{B}_n$ of length k?

```
• Artin length: coNP-complete(n, k) (Paterson & Razborov 1991); polynomial(n \le 3, k) (Sabalka 2003); open(n \ge 5, k);
```

How fast can you compute the complexity of a braid $\alpha \in \mathcal{B}_n$ of length k?

```
• Artin length:  \begin{array}{c} \text{coNP-complete}(n,k) & \text{(Paterson \& Razborov 1991);} \\ \text{polynomial}(n\leqslant 3,k) & \text{(Sabalka 2003);} \\ \text{open}(n\geqslant 5,k); & \end{array}
```

• Symmetric Garside length: polynomial (n, k) (Thurston 1988);

How fast can you compute the complexity of a braid $\alpha \in \mathcal{B}_n$ of length k?

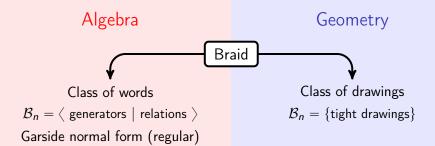
```
• Artin length: \text{coNP-complete}(n, k) (Paterson & Razborov 1991); \text{polynomial}(n \le 3, k) (Sabalka 2003); \text{open}(n \ge 5, k);
```

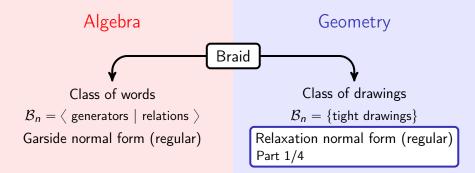
- Symmetric Garside length: polynomial (n, k) (Thurston 1988);
- Open laminated complexity: polynomial (n, k) (Dynnikov & Wiest 2004);

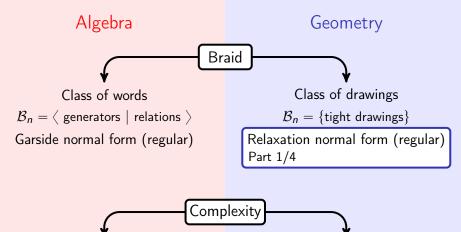
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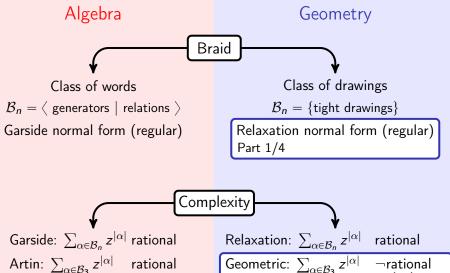
Garside: $\sum_{\alpha \in \mathcal{B}_n} z^{|\alpha|}$ rational

Artin: $\sum_{\alpha \in \mathcal{B}_3} z^{|\alpha|}$ rational

Artin: $\sum_{\alpha \in \mathcal{B}_{n \geq 4}} z^{|\alpha|}$?

rational

Relaxation: $\sum_{\alpha \in \mathcal{B}} z^{|\alpha|}$



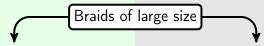
Artin: $\sum_{\alpha \in \mathcal{B}_{n \geqslant 4}} z^{|\alpha|}$?

Geometric: $\sum_{\alpha \in \mathcal{B}_3} z^{|\alpha|}$ —rational —algebraic —holonomic

Geometric: $\sum_{\alpha \in \mathcal{B}_{n > 4}} z^{|\alpha|}$?

Depth-first exploration

Width-first exploration



Random walk

Which normal forms converge? (Vershik, 2000)

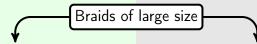
Markov-Ivanovsky normal form (Vershik & Malyutin, 2007)

Uniform measure on positive braids of given (Artin) size

What do Garside normal forms of large random braids look like?
(Gebhardt & Tawn, 2013)

Depth-first exploration

Width-first exploration



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Markov-Ivanovsky normal form (Vershik & Malyutin, 2007)

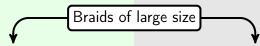
Garside normal forms
Part 3/4 (with J. Mairesse)

Uniform measure on positive braids of given (Artin) size

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What do Garside normal forms of large random braids look like? (Gebhardt & Tawn, 2013)

Uniform measure on infinite positive braids
Part 4/4 (with S. Abbes, S. Gouëzel & J. Mairesse)

Contents

- Geometric aspects of braids
 - Right relaxation normal form
 - Counting braids with a given geometric complexity
- 2 Algebraic aspects of braids
 - Garside normal form and random walks
 - Drawing infinite braids uniformly at random
- Conclusion

What is the right relaxation normal form?

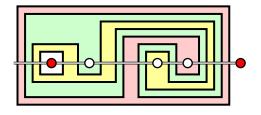
(by S. Caruso & B. Wiest)

Move your rightmost tensed puncture and relax!

What is the right relaxation normal form?

(by S. Caruso & B. Wiest)

Move your rightmost tensed puncture and relax!

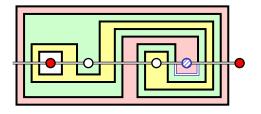


While your lamination is not trivial:

What is the right relaxation normal form?

(by S. Caruso & B. Wiest)

Move your rightmost tensed puncture and relax!

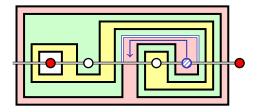


While your lamination is not trivial:

Select the rightmost (mobile) puncture that lies inside a bigon;

(by S. Caruso & B. Wiest)

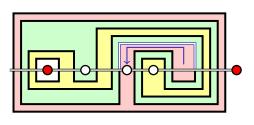
Move your rightmost tensed puncture and relax!



- Select the rightmost (mobile) puncture that lies inside a bigon;
- 2 Slide it along its right neighbour arc;

(by S. Caruso & B. Wiest)

Move your rightmost tensed puncture and relax!



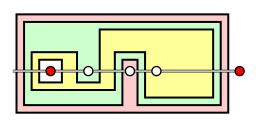
Moves performed:

[2 \(\sim 3 \)]

- Select the rightmost (mobile) puncture that lies inside a bigon;
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(by S. Caruso & B. Wiest)

Move your rightmost tensed puncture and relax!



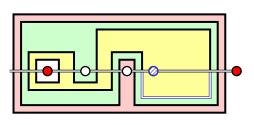
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(by S. Caruso & B. Wiest)

Move your rightmost tensed puncture and relax!



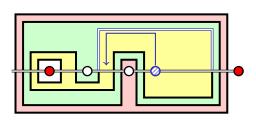
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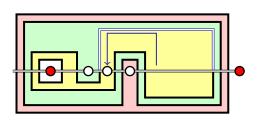
Move your rightmost tensed puncture and relax!



Moves performed:

[2 \(\sim 3]

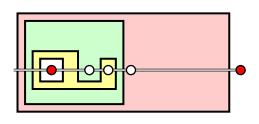
- Select the rightmost (mobile) puncture that lies inside a bigon;
- 2 Slide it along its right neighbour arc (and remember it);
- Relax your diagram!



Moves performed:

$$[2 \sim 3][2 \sim 3]$$

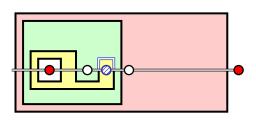
- Select the rightmost (mobile) puncture that lies inside a bigon;
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- Relax your diagram!



Moves performed:

$$[2 \sim 3][2 \sim 3]$$

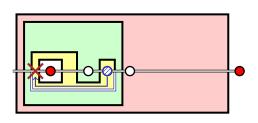
- Select the rightmost (mobile) puncture that lies inside a bigon;
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Moves performed:

$$[2 - 3][2 - 3]$$

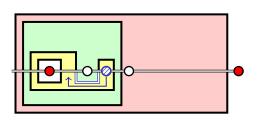
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- Relax your diagram!



Moves performed:

$$[2 \sim 3][2 \sim 3]$$

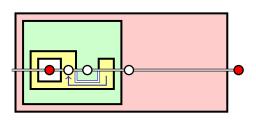
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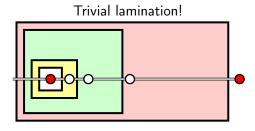
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- Relax your diagram!



Moves performed:

$$[2 - 3][2 - 3][1 - 2]$$

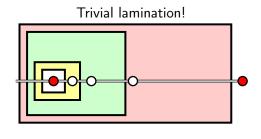
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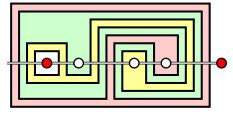
Relaxation normal form (RNF):

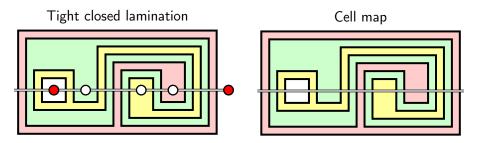
$$[1 \hookrightarrow 2] \cdot [2 \curvearrowright 3] \cdot [2 \curvearrowright 3]$$

$$\begin{bmatrix} k \smile \ell \end{bmatrix} = \sigma_k \dots \sigma_{\ell-1}$$
$$\begin{bmatrix} k \smile \ell \end{bmatrix} = \overline{\sigma_k} \dots \overline{\sigma_{\ell-1}}$$

- Select the rightmost (mobile) puncture that lies inside a bigon;
- Slide it along its right (or left) neighbour arc (and remember it);
- Relax your diagram!

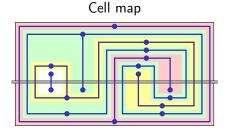
Tight closed lamination



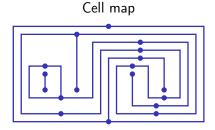


Tight closed lamination Cell map

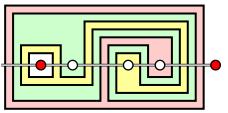
Tight closed lamination



Tight closed lamination

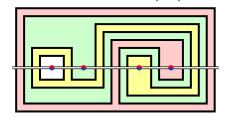


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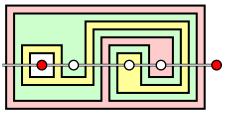


Cell map

Lamination trees (LT)

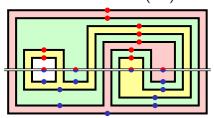


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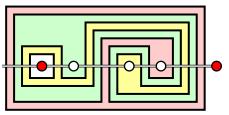


Cell map

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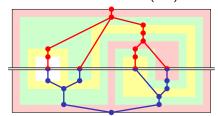


Tight closed lamination

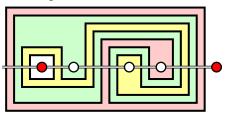


Cell map

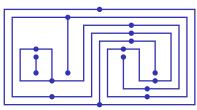
Lamination trees (LT)



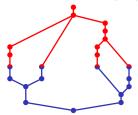
Tight closed lamination



Cell map



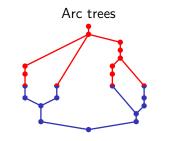
Lamination trees (LT)



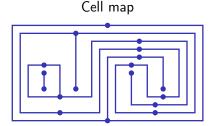
Tight closed lamination

Cell map

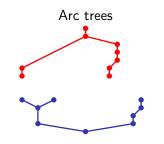
Lamination trees (LT)



Tight closed lamination

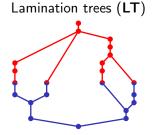


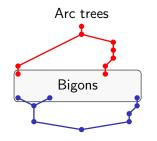
Lamination trees (LT)



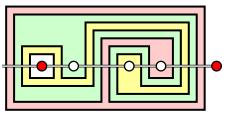
Tight closed lamination

Cell map



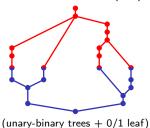


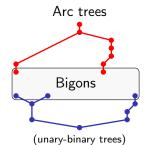
Tight closed lamination



Cell map

Lamination trees (LT)



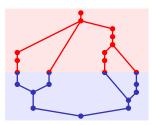


Given two braids $\alpha \in \mathcal{B}_n$ and $[k \curvearrowright \ell]$, remembering small-size subtrees $lt(\alpha)$ of $LT(\alpha)$ is enough to:

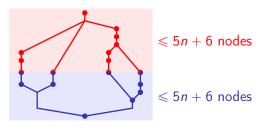
• check whether $RNF(\alpha[k \sim \ell]) = RNF(\alpha) \cdot [k \sim \ell];$

- check whether $RNF(\alpha[k \rightharpoonup \ell]) = RNF(\alpha) \cdot [k \rightharpoonup \ell];$
- ullet compute analogous subtrees $\mathbf{lt}(\alpha[k
 ightharpoonup \ell])$ (if needed).

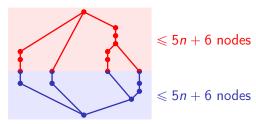
- check whether $\mathsf{RNF}(\alpha[k \rightharpoonup \ell]) = \mathsf{RNF}(\alpha) \cdot [k \rightharpoonup \ell];$
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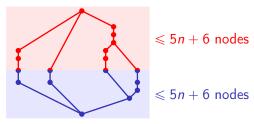


- check whether $\mathsf{RNF}(\alpha[k \rightharpoonup \ell]) = \mathsf{RNF}(\alpha) \cdot [k \rightharpoonup \ell];$
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Given two braids $\alpha \in \mathcal{B}_n$ and $[k \sim \ell]$, remembering small-size subtrees $lt(\alpha)$ of $LT(\alpha)$ is enough to:

- check whether $\mathsf{RNF}(\alpha[k \rightharpoonup \ell]) = \mathsf{RNF}(\alpha) \cdot [k \rightharpoonup \ell];$
- ullet compute analogous subtrees $\mathbf{lt}(\alpha[k
 ightharpoonup \ell])$ (if needed).



Theorem (J. 2015)

The right relaxation normal form is regular.

Going further

Some additional results

- Memory requirements: nearly optimal (up to a ratio ≤ 20);
- Dehornoy ordering: σ -positivity \Leftrightarrow **RNF** in a regular language.

Going further

Some additional results

- Memory requirements: nearly optimal (up to a ratio ≤ 20);
- Dehornoy ordering: σ -positivity \Leftrightarrow **RNF** in a regular language.

and open questions

- Is the right relaxation normal form (bi-)automatic? (Yes if $n \le 3$)
- Regularity of other transmission-relaxation normal forms? (wide open)

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Which complexity should we look at?

Knowing the open laminated complexity of α , can we compute its:

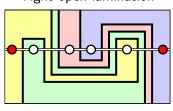
closed laminated complexity?

Knowing the open laminated complexity of α , can we compute its:

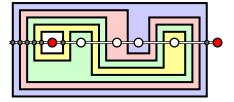
• closed laminated complexity?

Yes:
$$|\alpha|_c = |\alpha|_o + n + 3$$





Tight closed lamination



Knowing the open laminated complexity of α , can we compute its:

• closed laminated complexity?

Yes: $|\alpha|_c = |\alpha|_o + n + 3$

diagrammatic complexity?

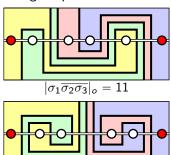
Knowing the open laminated complexity of α , can we compute its:

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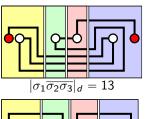
Yes:
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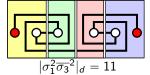
No for $n \geqslant 4$

Tight open laminations



Tight curve diagrams





 $|\sigma_1^2 \overline{\sigma_3}^2|_{\alpha} = 11$

Knowing the open laminated complexity of α , can we compute its:

• closed laminated complexity?

Yes: $|\alpha|_c = |\alpha|_o + n + 3$

diagrammatic complexity?

No for $n \ge 4$

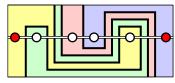
• inverse diagrammatic complexity?

Knowing the open laminated complexity of α , can we compute its:

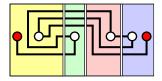
- closed laminated complexity?
- diagrammatic complexity?
- inverse diagrammatic complexity?

- $\text{Yes: } |\alpha|_c = |\alpha|_o + n + 3$
 - No for $n \ge 4$
 - Yes: $|\overline{\alpha}|_d = |\alpha|_o$

Tight open lamination



Inverse tight curve diagram



Knowing the open laminated complexity of α , can we compute its:

• closed laminated complexity?

Yes: $|\alpha|_c = |\alpha|_o + n + 3$

• diagrammatic complexity?

No for $n \ge 4$

• inverse diagrammatic complexity?

Yes: $|\overline{\alpha}|_d = |\alpha|_o$

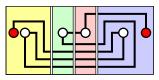
Let us compute geometric generating functions!

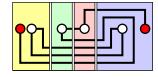
$$\begin{split} \mathcal{O}_n(z) &= \sum_{\alpha \in \mathcal{B}_n} z^{|\alpha|_o} \qquad \mathcal{C}_n(z) = \sum_{\alpha \in \mathcal{B}_n} z^{|\alpha|_c} \qquad \mathcal{D}_n(z) = \sum_{\alpha \in \mathcal{B}_n} z^{|\alpha|_d} \\ \mathcal{C}_n(z) &= z^{n+3} \mathcal{O}_n(z) = z^{n+3} \mathcal{D}_n(z) \end{split}$$

Generalising tight curve diagrams

Tight curve diagram

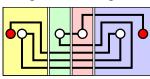
Tight generalised curve diagram



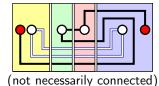


Generalising tight curve diagrams

Tight curve diagram

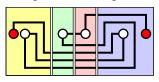


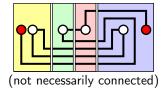
Tight generalised curve diagram



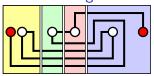
Generalising tight curve diagrams

Tight curve diagram Tight generalised curve diagram





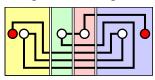
and encoding them!

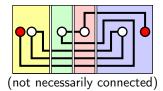


Coordinates: $\langle (x_0, x_1, x_2, x_3, x_4), (y_1, y_2, y_3, y_4) \rangle$

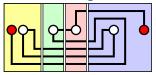
Generalising tight curve diagrams

Tight curve diagram Tight generalised curve diagram





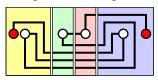
and encoding them!

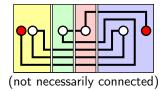


Coordinates: $\langle (0, x_1, x_2, x_3, 0), (y_1, y_2, y_3, y_4) \rangle$

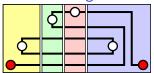
Generalising tight curve diagrams

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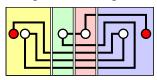
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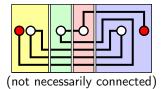


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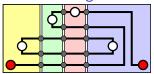
Generalising tight curve diagrams

Tight curve diagram Tight generalised curve diagram





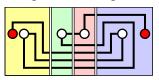
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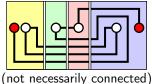


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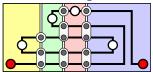
Generalising tight curve diagrams

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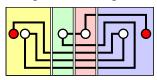
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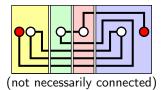


Coordinates: $\langle (0, 1, 2, 2, 0), (y_1, y_2, y_3, y_4) \rangle$

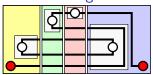
Generalising tight curve diagrams

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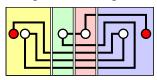
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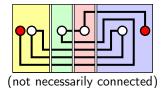


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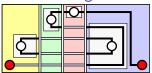
Generalising tight curve diagrams

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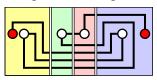
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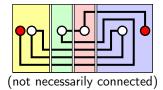


Coordinates: ((0,1,2,2,0),(1,3,4,0))

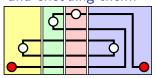
Generalising tight curve diagrams

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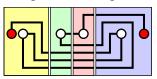
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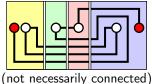


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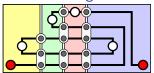
Generalising tight curve diagrams

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and encoding them!



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$$\textstyle \sum_{i=1}^{n-1} x_i = \frac{|\alpha|_d + 1 - n}{2} \text{:} \quad \text{ Let us compute } \mathcal{G}_n(z) = \sum_{k \geqslant 0} g_{n,k} z^k = \sqrt{z}^{1-n} \mathcal{D}_n(\sqrt{z})!$$

Theorem (J. 2015)

In the 3-strand braid group \mathcal{B}_3 , we have:

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$$\mathcal{G}_3(z)=2rac{1+2z-z^2}{z^2(1-z^2)}\left(\sum_{k\geqslant 3}\varphi(k)z^k
ight)+rac{1-3z^2}{1-z^2}$$
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$$\mathcal{G}_3(z)$$
 is more complicated than $\sum_{\alpha\in\mathcal{B}_3}z^{|\alpha|_{\mathsf{Artin}}}=\frac{(1+z)(1-z+z^2-2z^3)}{(2-z)(1-2z)(1-z-z^2)}$!

Contents

- Geometric aspects of braids
 - Right relaxation normal form
 - Counting braids with a given geometric complexity
- 2 Algebraic aspects of braids
 - Garside normal form and random walks
 - Drawing infinite braids uniformly at random
- Conclusion

Simple positive braids

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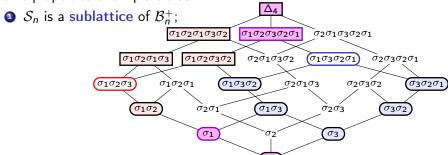
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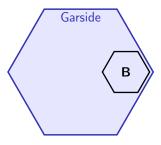
Consequence on Garside normal forms:

- **1** local neighbouring criterion: $w_1 \cdot w_2 \cdot \ldots \cdot w_k \in Gar(\mathcal{B}_n^+)$ iff
 - $w_1, \ldots, w_k \in \mathcal{S}_n \setminus \{1\};$ $\mathbf{R}(w_i) \supseteq \mathbf{L}(w_{i+1})$ for all i < k.

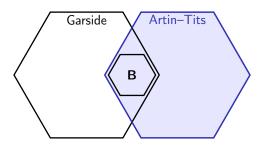
$$(\mathbf{L}(\alpha) = \{\sigma_i : \sigma_i \leqslant \alpha\} \text{ and } \mathbf{R}(\alpha) = \{\sigma_i : \alpha \geqslant \sigma_i\})$$



• Braid monoid: $\langle \sigma_i \mid \sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1}, i \neq j \pm 1 \Rightarrow \sigma_i \sigma_j = \sigma_j \sigma_i \rangle^+$;

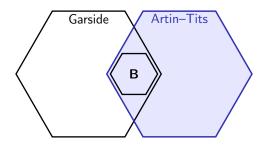


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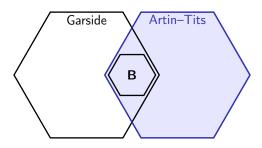
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$$\ell(i,j) = 2 \Rightarrow \sigma_i \sigma_j = \sigma_j \sigma_i$$

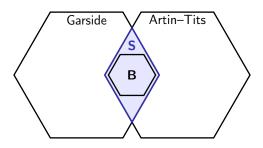


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$$\ell(i,j) = 3 \Rightarrow \sigma_i \sigma_j \sigma_i = \sigma_j \sigma_i \sigma_j$$

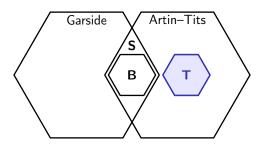


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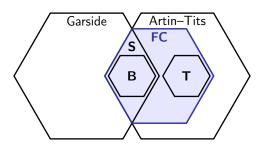
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Generalising braid monoids



- Braid monoid: $\langle \sigma_i \mid \sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1}, i \neq j \pm 1 \Rightarrow \sigma_i \sigma_j = \sigma_j \sigma_i \rangle^+$;
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- A–T monoid with spherical type: A–T and Garside; (3/4)
- Trace monoid: Artin–Tits with $\ell(i,j) \in \{2, +\infty\}$;
- A–T monoid with FC type: A–T with finite 2-way Garside family. (4/4)

Contents

- Geometric aspects of braids
 - Right relaxation normal form
 - Counting braids with a given geometric complexity
- 2 Algebraic aspects of braids
 - Garside normal form and random walks
 - Drawing infinite braids uniformly at random
- Conclusion

 $\bullet \ \, \text{Select i.i.d. generators } (Y_k)_{k\geqslant 0} \ \, \text{uniformly chosen in } \{\sigma_1,\ldots,\sigma_{n-1}\}.$

- **①** Select i.i.d. generators $(Y_k)_{k\geqslant 0}$ uniformly chosen in $\{\sigma_1,\ldots,\sigma_{n-1}\}$.
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① $Gar_{\ell}(\beta) = \beta_1 \cdot \ldots \cdot \beta_k$ with $R(\beta_i) \supseteq L(\beta_{i+1})$;

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Chronological **x** form:

Left Garside normal form:

1st step

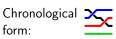
Right Garside \times normal form:

Left Garside X

- Select i.i.d. generators $(Y_k)_{k\geq 0}$ uniformly chosen in $\{\sigma_1,\ldots,\sigma_{n-1}\}$.
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Several Garside normal forms:

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Left Garside normal form:

2nd step

Right Garside normal form:

Left $^{\Delta}$ Garside \sim normal form:

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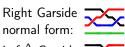
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Chronological form:



Left Garside normal form:





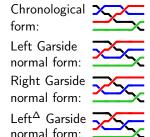
Left[△] Garside → normal form:

3rd step

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Several Garside normal forms:

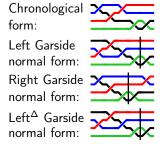
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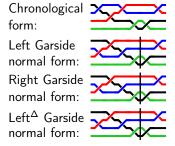
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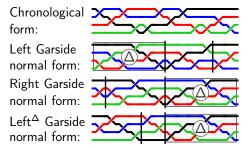
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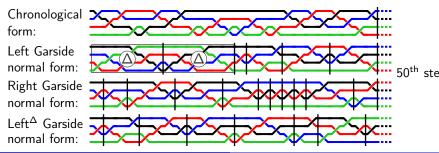
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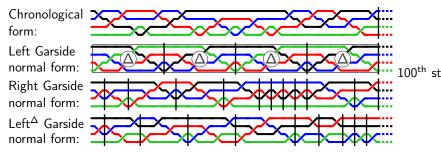
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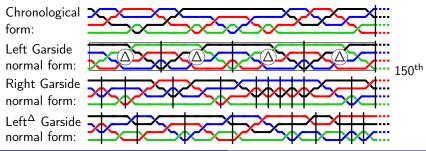
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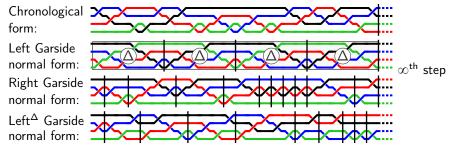
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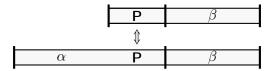
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Blocking pattern

Braid $P \in \mathcal{B}_n^+$ such that, for all $\alpha, \beta \in \mathcal{B}_n^+$ such that $\Delta_n \nleq \alpha P \beta$:

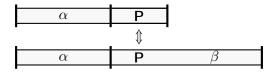
$$\textbf{0} \ \, \mathsf{Gar}_r(\mathsf{P}\beta) = \mathsf{Gar}_r(\mathsf{P}) \cdot \mathsf{Gar}_r(\beta) \ \, \mathsf{iff} \ \, \mathsf{Gar}_r(\alpha \mathsf{P}\beta) = \mathsf{Gar}_r(\alpha \mathsf{P}) \cdot \mathsf{Gar}_r(\beta);$$



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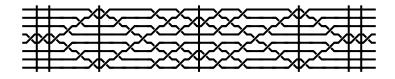
Blocking pattern

Braid $P \in \mathcal{B}_n^+$ such that, for all $\alpha, \beta \in \mathcal{B}_n^+$ such that $\Delta_n \nleq \alpha P \beta$:

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Blocking patterns exist in all braid monoids!

(Caruso & Wiest 2012)

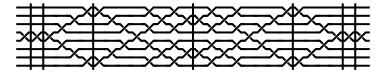


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Blocking patterns exist in all braid monoids! (Caruso & Wiest 2012) and in all irreducible A–T monoids of FC type!



Blocking pattern

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Some properties of blocking patterns

3 $\mathbf{C}_{\alpha\beta} \leqslant \mathbf{C}_{\alpha} + \mathbf{C}_{\beta} + \mathbf{K}$ for all braids $\alpha, \beta \in \mathcal{B}_{n}^{+}$ ($\mathbf{K} = \text{constant}$ and $\mathbf{C}_{x} = \#\{\text{occurrences of P or } \Delta_{n}^{-1}\mathbf{P}\Delta_{n} \text{ in } \mathbf{Gar}_{r}(x)\}$);

Blocking pattern

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Blocking patterns exist in all braid monoids! (Caruso & Wiest 2012) and in all irreducible A–T monoids of FC type!

Some properties of blocking patterns

- **3** $C_{\alpha\beta} \leq C_{\alpha} + C_{\beta} + K$ for all braids $\alpha, \beta \in \mathcal{B}_{n}^{+}$ ($K = \text{constant and } C_{x} = \#\{\text{occurrences of P or } \Delta_{n}^{-1}P\Delta_{n} \text{ in } Gar_{r}(x)\}\}$);

Blocking pattern

Braid $P \in \mathcal{B}_n^+$ such that, for all $\alpha, \beta \in \mathcal{B}_n^+$ such that $\Delta_n \nleq \alpha P \beta$:

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Blocking patterns exist in all braid monoids! (Caruso & Wiest 2012) and in all irreducible A–T monoids of FC type!

Some properties of blocking patterns

- $\mathbf{S} \ \, \mathbf{C}_{\alpha\beta} \leqslant \mathbf{C}_{\alpha} + \mathbf{C}_{\beta} + \mathbf{K} \ \, \text{for all braids} \ \, \alpha, \beta \in \mathcal{B}_{n}^{+} \\ \ \, (\mathbf{K} = \text{constant and} \ \, \mathbf{C}_{x} = \#\{\text{occurrences of } \mathbf{P} \ \, \text{or} \ \, \Delta_{n}^{-1} \mathbf{P} \Delta_{n} \ \, \text{in} \ \, \mathbf{Gar}_{r}(x)\});$

Theorem (J. & Mairesse 2016⁺)

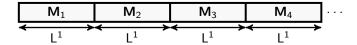
Prefixes of the words $Gar_r(X_k)_{k\geq 0}$ almost surely converge.





What is our limit object? How fast do we reach it?

Limit of an infinite-state Markov chain with L¹ factors;





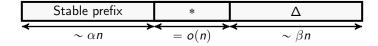
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Going even further

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Convergence with arbitrarily large steps?

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(wide open)

Contents

- Geometric aspects of braids
 - Right relaxation normal form
 - Counting braids with a given geometric complexity
- 2 Algebraic aspects of braids
 - Garside normal form and random walks
 - Drawing infinite braids uniformly at random
- Conclusion



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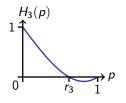
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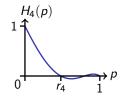
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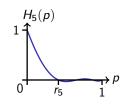
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What if
$$H_n(p) \rightarrow 0$$
? (i.e. $p \rightarrow r_n$)







Two-step limit extraction

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Theorem (Abbes, Gouëzel, J. & Mairesse 2016+)

Uniform probability measures on \mathcal{B}_n^k converge weakly towards μ_∞ when $k \to +\infty$.

 μ_{∞} is a uniform probability measure on infinite braids!

Stable region conjectures (Gebhardt & Tawn 2014)

1 The words $\{\operatorname{Gar}_{\ell}(\beta) \mid \beta \in \partial \mathcal{B}_n^+\}$ contain a geometric number of Δ_n ;

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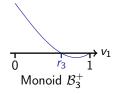
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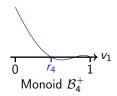
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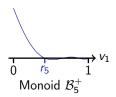
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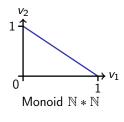


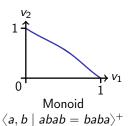
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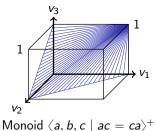
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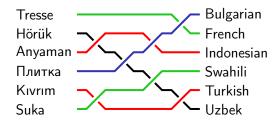
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Uniform measures

 Uniform measures on positive spheres converge towards a simple critical Markov process (for all irreducible A–T monoids of FC type).



Do you have questions?

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