# Functional programming Lecture 02 — First steps

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# Schedule

- 1. Types & functions
- 2. Function on lists
- 3. High-order functions & list comprehensions
- 4. Test
- 5. Algebraic data type Binary search trees
- 6. Algebraic data type General trees
- 7. Problem solving Countdown
- 8. Functional images
- 9. File system
- 10. Radar or Cloud simulator
- 11. Test

# **Functional programming concepts**

Functional programming concepts

First steps

Types and classes

Defining functions

Some functions

# Genealogy of programming languages



## **Functional languages**















## Lisp

Lisp (historically, LISP) is a family of computer programming languages with a long history and a distinctive, fully parenthesized prefix notation. Originally specified in 1958, Lisp is the second-oldest high-level programming language in widespread use today. (Only Fortran is older, by one year.)



## Erlang

Erlang is a general-purpose, concurrent, functional programming language, as well as a garbage-collected runtime system.



## **Elixir** Elixir is a functional, concurrent, general-purpose programming language that runs on the Erlang virtual machine (BEAM).



F# F# is a strongly typed, multi-paradigm programming language that encompasses functional, imperative, and object-oriented programming methods. It is being developed at Microsoft Developer Division and is being distributed as a fully supported language in the .NET framework.

# Caml

#### Ocaml

Ocaml, originally named Objective Caml, is the main implementation of the programming language Caml. OCaml's toolset includes an interactive top-level interpreter, a bytecode compiler, a reversible debugger, a package manager (OPAM), and an optimizing native code compiler.



#### **Clojure**

Clojure is a dialect of the Lisp programming language. Clojure is a general-purpose programming language with an emphasis on functional programming. It runs on the Java virtual machine and the Common Language Runtime.



#### Racket

Racket, formerly PLT Scheme, is a general purpose,

multi-paradigm programming language in the Lisp-Scheme family.

One of its design goals is to serve as a platform for language

creation, design, and implementation



#### Elm

Elm is a domain-specific programming language for declaratively creating web browser-based graphical user interfaces. Elm is purely functional, and is developed with emphasis on usability, performance, and robustness.



# Scala Scala is a general-purpose programming language providing support for functional programming and a strong static type system. Designed to be concise, many of Scala's design decisions aimed to address criticisms of Java.



#### Haskell

Haskell is a general-purpose, statically-typed, purely functional programming language with type inference and lazy evaluation. Designed for teaching, research and industrial applications, Haskell has pioneered a number of programming language features such as type classes, which enable type-safe operator overloading, and monadic IO. Haskell's main implementation is the Glasgow Haskell Compiler (GHC). It is named after logician Haskell Curry.

## Characteristics of functional programming



# Haskell

- Haskell is a compiled, statically typed, functional programming language.
- It was created in the early 1990s as one of the first open-source purely functional programming languages.
- It is named after the American logician Haskell Brooks Curry.



# Glasgow Haskell Compiler

- Concise programs
- Powerful type system
- List comprehensions
- Recursive functions
- High-order functions
- Effectful functions
- Generic functions
- Lazy evaluation
- Equational reasoning

## The imperatives

- GHC: state-of-the-art, open source, compiler and interactive environment for the functional language Haskell.
- GHCi: GHC's interactive environment.
- Hackage: Haskell community's central package archive of open source software.

## **Testing Frameworks**

- QuickCheck: powerful testing framework where test cases are generated according to specific properties.
- HUnit: unit testing framework similar to JUnit.
- Hspec: a testing framework similar to RSpec with support for QuickCheck and HUnit.
- The Haskell Test Framework, HTF: integrates both Hunit and QuickCheck.

## **Ancillary Tools**

- darcs: revision control system.
- haddock: documentation system.
- cabal: build system.
- stack: build system.
- hoogle: type-aware API search engine.

## **Static Analysis Tools**

- hlint: detect common style mistakes and redundant parts of syntax, improving code quality.
- Sourcegraph: Haskell visualizer.

## **Dynamic Analysis Tools**

- criterion: powerful benchmarking framework.
- hpc: check evaluation coverage of a haskell program, useful for determining test coverage.

# IDEs

- VSCodium.
- IntelliJ.
- Vim.
- GNU Emacs.
- Haskell for Mac (commercial).
- Sublime Text (commercial)



## Haskell books



## Functional programming books





sum :: Num a => [a] -> a
sum [] = 0
sum (x : xs) = x + sum xs

```
sum :: Num a => [a] -> a
sum [] = 0
sum (x : xs) = x + sum xs
 sum [1,2,3]
= { applying function sum }
 1 + sum [2,3]
= { applying function sum }
 1 + (2 + sum [3])
= { applying function sum }
1 + 2 + (3 + sum ])
= { applying function sum }
1 + (2 + (3 + 0))
= { applying function + }
 1 + (2 + 3)
= { applying function + }
1 + 5
= { applying function + }
 6
```

```
qsort :: Ord a => [a] -> [a]
qsort [] = []
qsort (x : xs) = qsort smaller ++ [x] ++ qsort larger
where
smaller = [x' | x' <- xs, x' <= x]
larger = [x' | x' <- xs, x' > x]
```

= { applying function qsort }

= { applying function ++ }

[] ++ [x] ++ []

[x]

```
qsort :: Ord a => [a] -> [a]
qsort [] = []
qsort (x : xs) = qsort smaller ++ [x] ++ qsort larger
where
    smaller = [x' | x' <- xs, x' <= x]
    larger = [x' | x' <- xs, x' > x]
qsort [x]
= { applying function qsort }
qsort [] ++ [x] ++ qsort [x]
```

```
qsort :: Ord a \Rightarrow [a] \rightarrow [a]
gsort [] = []
qsort (x : xs) = qsort smaller ++ [x] ++ qsort larger
  where
     smaller = [x' | x' < xs, x' < x]
    larger = [x' | x' < xs, x' > x]
 qsort [3,5,1,4,2]
= { applying function qsort }
 qsort [1,2] ++ [3] ++ qsort [5,4]
= { applying function qsort }
 (qsort [] ++ [1] ++ qsort [2]) ++ [3] ++ (qsort [4] ++ [5] ++ qsort [])
= { applying function qsort }
 ([] ++ [1] ++ [2]) ++ [3] ++ ([4] ++ [5] ++ [])
= { applying function ++ }
 [1,2] ++ [3] ++ [4,5]
= { applying function ++ }
 [1, 2, 3, 4, 5]
```



Functional programming concepts

## First steps

Types and classes

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Some functions

- The *Glasgow Haskell Compiler* (GHC) is the state-of-the-art open source implementation of Haskell
- The GHC if freely available for a range of operating systems from the Haskell home page http://www.haskell.org
- We recommand downloading the Haskell Platform
- Once installed, the interface GHCi system can be started from the terminal command prompt by simply typing ghci.

# GHCi

 $\lambda > 1+2+3$ 6  $\lambda > 1+2*3$ 7  $\lambda > (1+2)*3$ 9  $\lambda > 2-3+4$ 3  $\lambda > 2*3/4$ 1.5

λ > 2\*pi
6.283185307179586

 $\label{eq:lambda} \begin{array}{l} \lambda > \mbox{(1 + sqrt 5) / 2} \\ 1.618033988749895 \end{array}$ 

 $\label{eq:log2} \begin{array}{l} \lambda > \mbox{log2} \\ 0.6931471805599453 \end{array}$ 

# GHCi

```
\begin{array}{ll} \lambda > & 2^3 \mbox{$^2$} \\ 2417851639229258349412352 \end{array}
```

```
\lambda > (2^3)^4
4096
\lambda > ceiling 2.6
3
\lambda > floor 2.6
2
\lambda > round 2.6
3
\lambda > (sin pi)<sup>2</sup> + (cos pi)<sup>2</sup>
1.0
```
$\begin{aligned} \lambda > & x = 42 \\ \lambda > & x+1 \\ 43 \end{aligned}$ 

 $\lambda$  > let x = 42 in x+1 43

```
\label{eq:laskell} \begin{array}{ll} \lambda > & "\text{Haskell rocks!"} \\ & "\text{Haskell rocks!"} \end{array}
```

```
\lambda > "Haskell " ++ "rocks!"
"Haskell rocks!"
```

```
λ > "Haskell " <> "rocks!"
"Haskell rocks!"
```

$$\label{eq:linear} \begin{split} \lambda > & ['H', 'a', 's', 'k', 'e', 'l', 'l', '', 'r', 'o', 'c', 'k', 's', '!'] \\ "Haskell rocks!" \end{split}$$

Command	Meaning
:load name	load script name
:reload	reload current script
:set editor name	set editor to <i>name</i>
:edit name	edit script name
:edit	edit current script
:type <i>expr</i>	show type of <i>expr</i>
:?	show all commands
:quit	quit GHCi

 $\lambda$  > :type 1 1 :: Num a => a

 $\lambda$  > :type 2.5 2.5 :: Fractional a => a

 $\lambda$  > :type 5/2 5/2 :: Fractional a => a

λ > :type 5 `div` 2
5 `div` 2 :: Integral a => a

```
\lambda > :type 1+2
1+2 :: Num a => a
\lambda > :type (+)
(+) :: Num a => a -> a -> a
\lambda > :type (1 +)
(1 +) :: Num a => a -> a
\lambda > : type (+ 1)
(+ 1) :: Num a => a -> a
```

 $\lambda$  > :type 2.5 2.5 :: Fractional a => a  $\lambda$  > :type 5/2 5/2 :: Fractional a => a

λ > :type (/)
(/) :: Fractional a => a -> a -> a

 $\lambda$  > :type (/ 2) (/ 2) :: Fractional a => a -> a

```
\lambda > :type pi
pi :: Floating a => a
```

```
\lambda > :type sqrt 2
sqrt 2 :: Floating a => a
```

```
\lambda > :type cos
cos :: Floating a => a -> a
```

### GHCi

```
\lambda > fact n = if n == 0 then 1 else n * fact (n-1)
```

```
\lambda > :type fact
fact :: (Eq a, Num a) \Rightarrow a \Rightarrow a
\lambda > fact 5
120
\lambda > fact 0
1
\lambda > fact 5.0
120.0
\lambda > fact 2.5
<sup>CInterrupted.</sup>
```

 $\lambda >$  f = fact

 $\lambda>$  :type fact fact :: (Eq a, Num a) => a -> a

 $\lambda >$  f 5 120

 $\lambda >$  f (f 3) 720  $\lambda >$  'a' 'a'

 $\lambda >$  'abc' <interactive>: error: o Syntax error on 'abc'  $\lambda >$  'a':"bc"

"abc"

 $\lambda >$  "abc" "abc"

 $\lambda$  > :type "abc" "abc" :: String

 $\lambda > \text{ "abc" ++ "def"} \\ \text{"abcdef"}$ 

λ > :type (++) (++) :: [a] -> [a] -> [a]

# Types and classes

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- In Haskell every expression must have a type.
- A type is a collection of related values.
- We use the notation  $v \ :: \ T$  to mean that v is a value in the type T.

#### Example

True	:: Bool
False	:: Bool
not	:: Bool -> Bool
(&&)	:: Bool -> Bool -> Bool
(  )	:: Bool -> Bool -> Bool

- Bool Logical values.
- Char Single characters.
- String Strings of characters.
- Int Fixed-precision integers.
- Integers Arbitrary-precision integers.
- Float Since-precision floating-point numbers.
- Double Double-precision floating-point numbers.

- A list is a sequence of elements of the same type, with the elements being enclosed in square parentheses and separated by commas.
- We write [T] for the type of all lists whose elements have type T.
- The number of elements in a list is called its length.
- The list [] of length zero is called the empty list.
- [] and [[]] (and [[[]]], [[[[]]], ...) are different lists.

λ > :type [] [] :: [a]

λ > :type [1,2,3,4,5] [1,2,3,4,5] :: Num a => [a]

$$\label{eq:lassian} \begin{split} \lambda > \ : type \ ['a', 'b', 'c', 'd'] \\ ['a', 'b', 'c', 'd'] \ :: \ [Char] \end{split}$$

$$\label{eq:label} \begin{split} \lambda > \ : type \ ["ab", "cd", "ef", "gh"] \\ ["ab", "cd", "ef", "gh"] \ :: \ [String] \end{split}$$

```
λ > :type [cos, sin]
[cos, sin] :: Floating a => [a -> a]
λ > :type [1, 'a']
<interactive>: error:
    o No instance for (Num Char) arising from the literal '1'
λ > :type [[1],[2,3],[4,5,6]]
[[1],[2,3],[4,5,6]] :: Num a => [[a]]
```

 $\lambda > :type [[[1]], [[2,3], [4,5,6]]]$ [[[1]], [[2,3], [4,5,6]]] :: Num a => [[[a]]]

- A tuple is a sequence of components of possibly different types, with the components being enclosed in round parentheses and separated by commas.
- We write (T1, T2, ..., Tn) for the type of all tuples whose *i*-th component have type Ti for any 1 ≤ *i* ≤ *n*.
- The number of elements in a tuple is called its arity.
- The tuple () of arity zero is called the empty tuple.
- Tuple of arity one are not permitted.

 $\lambda$  > :type () () :: ()

```
\lambda > :type (1,'a')
(1,'a') :: Num a => (a, Char)
```

```
λ > :type (1,2,'a',"abc")
(1,2,'a',"abc") :: (Num a, Num b) => (a, b, Char, String)
```

```
\lambda > :type (sqrt, 'a')
(sqrt, 'a') :: Floating a => (a -> a, Char)
```

```
λ > :type (1, ('a', "cd"))
(1, ('a', "cd")) :: Num a => (a, (Char, String))
```

#### Tuple types

```
λ > :type (1, ('a', "cd"))
(1, ('a', "cd")) :: Num a => (a, (Char, String))
```

```
λ > :type (1, [cos, sin])
(1, [cos, sin]) :: (Floating a1, Num a2) => (a2, [a1 -> a1])
```

```
\lambda > :type (1)
(1) :: Num a => a
```

 $\lambda$  > let t = (1,2) in (t, 3) ((1,2),3)

 $\lambda$  > let t = (1,t)

<interactive>: error:

o Couldn't match expected type 'b' with actual type '(a, b)'

- A function is a mapping of one type to results of another type.
- We write T1 -> T2 for the type of all functions that map arguments of type T1 to results of type T2.
- There is no restriction that function must be total on their argument type.

#### **Function types**

```
\lambda > :type not
not :: Bool -> Bool
```

```
\lambda > :type even -- :type odd
even :: Integral a => a -> Bool
```

```
\lambda > :type mod
mod :: Integral a => a -> a -> a
```

```
\lambda > add x y = x+y
\lambda > :type add
add :: Num a => a -> a -> a
```

```
\lambda > add' (x,y) = x+y
\lambda > :type add'
add' :: Num a => (a, a) -> a
```

- Currying is the process of transforming a function that takes multiple arguments in a tuple as its argument, into a function that takes just a single argument and returns another function which accepts further arguments, one by one, that the original function would receive in the rest of that tuple.
- The function arrow -> in type is assumed to associate to the right.

The type T1 -> T2 -> T3 -> ... -> Tn means T1 -> (T2 -> (T3 -> ( ... -> Tn)...)) The type a1 -> a2 -> a3 means a1 -> (a2 -> a3) The type a1 -> a2 -> a3 -> a4 means a1 -> (a2 -> (a3 -> a4)) The type a1 -> a2 -> a3 -> a4 -> a5 means a1 -> (a2 -> (a3 -> (a4 -> a5)))

Multiplying three integers

-- mult :: Int -> (Int -> (Int -> Int)) mult :: Int -> Int -> Int -> Int mult x y z = x\*y\*z

```
Multiplying three integers
-- mult :: Int \rightarrow (Int \rightarrow (Int \rightarrow Int))
mult :: Int -> Int -> Int -> Int
mult x y z = x*y*z
\lambda > mult 2 3 4
24
\lambda > :type mult 2
mult 2 :: Int -> Int -> Int
\lambda > :type mult 2 3
mult 2 3 :: Int -> Int
\lambda > :type mult 2 3 4
mult 2 3 4 :: Int
```

Multiplying three integers

```
-- mult :: Int -> (Int -> (Int -> Int))

mult :: Int -> Int -> Int -> Int

mult x y z = x*y*z

λ > mult2 = mult 2

λ > mult3 = mult2 3

λ > mult3 4

24
```

```
\lambda > :type mult2
mult2 :: Int -> Int -> Int
\lambda > :type mult3
mult3 :: Int -> Int
```

- Parametric polymorphism refers to when the type of a value contains one or more (unconstrained) type variables, so that the value may adopt any type that results from substituting those variables with concrete types.
- For example, the function id :: a -> a contains an unconstrained type variable a in its type, and so can be used in a context requiring Char -> Char or Integer -> Integer or (Bool -> Bool) -> (Bool -> Bool) or any of a literally infinite list of other possibilities.
- The empty list [] :: [a] belongs to every list type.

```
\lambda > length []
0
\lambda > length [1,3,5,7,2,4,6,8]
8
\lambda > length ["Huey", "Dewey", "Louie"]
3
\lambda > length [sin, cos, tan]
3
```

```
\lambda > :type length
length :: Foldable t => t a -> Int
```

```
λ > :info length
type Foldable :: (* -> *) -> Constraint
class Foldable t where
  length :: t a -> Int
   ...
   -- Defined in 'Data.Foldable'
```

- A type that contains one or more class constraints is called overloaded.
- Class constraints are written in the form C a, where C is the name of the class and a is a type variable.

#### **Overloaded types**

 $\lambda > 1 + 2$ 3  $\lambda$  > :type 1 1 :: Num a  $\Rightarrow$  a  $\lambda$  > :type 1 + 2 1 + 2 :: Num a => a $\lambda$  > sqrt 2 + sqrt 3 3.1462643699419726  $\lambda$  > :type sqrt 2 sqrt 2 :: Floating a => a  $\lambda$  > :type sqrt 2 + sqrt 3 sqrt 2 + sqrt 3 :: Floating a => a

 $\lambda > 1.0 + 2.0$ 

3.0

 $\lambda >$  :type 1.0

1.0 :: Fractional a => a

 $\lambda$  > :type 1.0 + 2.0 1.0 + 2.0 :: Fractional a => a

#### **Overloaded types**

```
\lambda > :type sqrt
sqrt :: Floating a => a -> a
```

- A class is collection of types that support certain overloaded operations called methods.
- Haskell provides a number of basic classes that are built-in to the language.
# Eq – Equality types

This class contains types whose values can be compared for equality and inequality using the following two methods:

(==) :: a -> a -> Bool (/=) :: a -> a -> Bool

All the basic types Bool, Char, String, Int, Integers, Float and Double are instances of the Eq class.

Eq - Equality types $\lambda > True == True$ True

 $\label{eq:lambda} \lambda > \texttt{'a'} == \texttt{'b'}$  False

 $\lambda$  > "abc" == "abc"

True

```
Eq – Equality types
\lambda > ('a', 1) == ('b', 1)
False
\lambda > (1, 2, 3) == (1, 2)
<interactive>:120:14: error:
  o Couldn't match expected type: (a0, b0, c0)
    with actual type: (a1, b1)
\lambda > [1,2,3] == [1,2,3,4]
False
\lambda > cos == cos
<interactive>: error:
  o No instance for (Eq (Double -> Double))
    arising from a use of '=='
```

# Ord - Ordered types

This class contains types that are instances of the equality class Eq, but in addition these values are totally ordered, and as such can be compared using the following six methods:

(<) :: a -> a -> Bool (<=) :: a -> a -> Bool (>) :: a -> a -> Bool (>=) :: a -> a -> Bool min :: a -> a -> Bool max :: a -> a -> a

All the basic types Bool, Char, String, Int, Integers, Float and Double are instances of the Ord class.

```
\begin{array}{ll} \mbox{Ord} - \mbox{Ordered types} \\ \lambda > & \mbox{False} < \mbox{True} \\ \mbox{True} \end{array}
```

 $\lambda > \text{ "elegant" < "elephant"}$  True

```
\lambda > \text{ "a" < "ab"}
```

True

 $\lambda > \ \texttt{'b'} > \ \texttt{'a'}$ 

True

```
\lambda > [1,2,3] <= [1,2]
False
```

```
\begin{array}{ll} \lambda > & [] & < & [1] \\ True \end{array}
```

# $\label{eq:constraint} \begin{array}{l} \mbox{Ord} - \mbox{Ordered types} \\ \lambda > \ (1,2) \ < \ (1,3) \\ \mbox{True} \end{array}$

 $\lambda >$  (1,2,3) < (1,1)

<interactive>: error:

o Couldn't match expected type: (a0, b0, c0)
with actual type: (a1, b1)

```
\label{eq:linear} \begin{array}{ll} \lambda > & [\texttt{True}] < [\texttt{False},\texttt{False}] \\ \texttt{False} \end{array}
```

```
\lambda > (False,False) <= (False,True) 
 True
```

```
Ord – Ordered types
\lambda >
\lambda > \min('a', 2)('a', 1)
('a',1)
\lambda > \max('a', 2)('a', 1)
('a',2)
\lambda > \sin < \cos
<interactive>: error:
  o No instance for (Ord (Double -> Double))
    arising from a use of '<'
\lambda > (1, \sin) > (2, \cos)
<interactive>: error:
  o No instance for (Ord (Double -> Double))
    arising from a use of '>'
```

#### Show – Showable types

This class contains types that can be converted into strings of characters using the following method:

show :: a -> String

All the basic types Bool, Char, String, Int, Integers, Float and Double are instances of the Show class.

```
Show – Showable types
\lambda > show True
"True"
\lambda > show 'a'
"'a'"
\lambda > show "abc"
"\"abc\""
\lambda > show [1,2,3]
"[1,2,3]"
\lambda > show (1, True, [1,2,3])
```

"(1,True,[1,2,3])"

#### Read – Readable types

This class is dual to **Read** and contains types whose values can be converted from string of characters using the following method:

```
read :: String -> a
```

All the basic types Bool, Char, String, Int, Integers, Float and Double are instances of the Read class.

```
Read – Readable types
\lambda > read "False" :: Bool
False
\lambda > read "'a'" :: Char
'a'
\lambda > read "\"abc\"" :: String
"abc"
\lambda > read "[1,2,3]" :: [Int]
[1, 2, 3]
\lambda > read "(1, True, [1,2,3])" :: (Int, Bool, [Int])
(1,True, [1,2,3])
```

#### Num – Numeric types

This class contains types whose values are numeric, and as such can be processed using the following six methods:

(+) :: a -> a -> a (-) :: a -> a -> a (\*) :: a -> a -> a negate :: a -> a abs :: a -> a signum :: a -> a

Note that the Num class does not provide a division method.

Num	<ul> <li>Numeric types</li> </ul>
<mark>λ</mark> > 3	1+2
<b>λ</b> > -1	1-2
<mark>λ</mark> > 3.0	1.0+2.0
<mark>λ</mark> > 6	2*3
λ > 6.0	2.0*3.0

```
Num – Numeric types
\lambda > negate 3.0
-3.0
\lambda > negate (-2)
2
\lambda > abs(-1.5)
1.5
\lambda > signum 3
1
\lambda > signum (-3)
-1
```

#### Integral – Integral types

This class contains types that are instances of the numeric class Num, but in addition whose values are integers, and as such support the method of integer division and integer remainder:

div :: a -> a -> a mod :: a -> a -> a

```
Integral – Integral types
\lambda > \text{div } 7 2
3
\lambda > 7 'div' 2
3
\lambda > 8 'div' 2
4
\lambda > 7 \mod 2
1
\lambda > 8 \mod 2
0
```

```
Integral – Integral types
\lambda > (-7) `div` 2
-4
\lambda > (-7) 'div' (-2)
3
\lambda > (-7) \mod 2
1
\lambda > (-7) `mod` (-2)
-1
```

#### Fractional - Fractional types

This class contains types that are instances of the numeric class Num, but in addition whose values are non-integral, and as such support the method of integer fractional division and fractional reciprocation:

(/) :: a -> a -> a recip :: a -> a -> a

The basic types Float and Double are instances of the Fractional class.

#### Fractional - Fractional types

 $\lambda > 7.0 / 2.0$ 3.5

 $\begin{array}{l} \lambda > \ 2.0 \ / \ 7.0 \\ 0.2857142857142857 \end{array}$ 

 $\begin{array}{ll} \lambda > \text{ recip } 2.0 \\ 0.5 \end{array}$ 

λ> recip 1.0 1.0

# **Defining functions**

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Some functions

```
even :: Integral a => a -> Bool
even n = n `mod` 2 == 0
```

```
odd :: Integral a => a -> Bool
odd n = n `mod` 2 /= 0
```

```
recip :: Fractional a => a -> a
recip n = 1 / n
```

For processing conditions, the if-then-else syntax was defined in Haskell98.

if <condition> then <true-value> else <false-value>

if is an expression (which is converted to a value) and not a statement (which is executed) as in many imperative languages. As a consequence, the else is mandatory in Haskell. Since if is an expression, it must evaluate to a result whether the condition is true or false, and the else ensures this.

# **Conditional expressions**

```
abs :: Int \rightarrow Int
abs n = if n >= 0 then n else -n
```

```
signum :: Int -> Int
signum n = if n < 0 then -1 else
if n == 0 then 0 else 1</pre>
```

```
describeLetter :: Char -> String
describeLetter c = if c >= 'a' && c <= 'z'
            then "Lower case"
            else if c >= 'A' && c <= 'Z'
            then "Upper case"
            else "Not an ASCII letter"</pre>
```

addOneIfEven1 :: Integral a => a -> a addOneIfEven1 n = if even n then n+1 else n

addOneIfEven2 :: Integral a => a -> a addOneIfEven2 n = n + if even n then 1 else 0

addOneIfEven3 :: Integral a => a -> a addOneIfEven3 n = (if even n then (+ 1) else (+ 0)) n

addOneIfEven4 :: Integral a => a -> a addOneIfEven4 n = (if even n then (+ 1) else id) n Remember that

```
isNullLength :: Foldable t => t a -> Bool
isNullLength xs = if length xs == 0 then True else False
is nothing but
isNullLength :: Foldable t => t a -> Bool
isNullLength xs = length xs == 0
or (as we we shall see soon ... but not really better here!)
isNullLength :: Foldable t => t a -> Bool
isNullLength = (== 0) . length
```

As an alternative to using conditional expressions, functions can also be defined using guarded expressions, in which a sequence of logical expressions called guards is used to choose between a sequence of results of the same type.

- If the first guard is **True**, then the first result is chosen.
- Otherwise, if the second guard is **True**, then the second result is chosen.
- And so on.

abs1 :: Int  $\rightarrow$  Int abs1 n = if n >= 0 then n else -n

abs2 :: Int -> Int abs2 n

| n >= 0 = n

| otherwise = -n

# **Guarded expressions**

```
signum2 :: Int -> Int
signum2 n
| n < 0 = -1
| n == 0 = 0
| otherwise = 1</pre>
```

# **Guarded expressions**

describeLetter2 :: Char -> String describeLetter2 c | c >= 'a' && c <= 'z' = "Lower case" | c >= 'A' && c <= 'Z' = "Upper case" | otherwise = "Not an ASCII letter" fact :: (Eq t, Num t)  $\Rightarrow$  t  $\rightarrow$  t fact n | n == 0 = 1| otherwise = n \* fact (n-1)mult :: (Eq t, Num t, Num a)  $\Rightarrow$  a  $\rightarrow$  t  $\Rightarrow$  a mult n m | m == 0 = 0| otherwise = n + mult n (m - 1)

Many functions have a simple and intuitive definition using pattern matching, in which a sequence of syntactic expressions called patterns is used to choose between a sequence of results of the same type.

The wildcard pattern \_ matches any value.

- If the first pattern is matched, then the first result is chosen.
- Otherwise, if the second pattern is matched, then the second result is chosen.
- And so on...

# Pattern matching

-- conditional expression
not :: Bool -> Bool
not b = if b == True then False else True

```
-- guarded function
not :: Bool -> Bool
not b
| b == True = False
| otherwise = True
-- pattern matching
not :: Bool -> Bool
```

not False = True

not True = False

# Pattern matching

(&&) :: Bool -> Bool -> Bool True && True = True True && False = False False && True = False False && False = False

```
(&&) :: Bool -> Bool -> Bool
True && b = b
False && _ = False
```

### Pattern matching

guess	:	: 1	Int	->	String
guess	0	=	"I	am	zero"
guess	1	=	"I	am	one"
guess	2	=	"I	am	two"
guess	_	=	"I	am	at least three"

```
-- be careful with the wildcard pattern
guess :: Int -> String
guess _ = "I am at least three"
guess 0 = "I am zero"
guess 1 = "I am one"
guess 2 = "I am two"
```

A tuple of patterns is itself a pattern, which matches any tuple of the same arity whose components all match the corresponding patterns in order.

first :: (a, b, c) -> a
first (x, \_, \_) = x

second :: (a, b, c) -> b
second (\_, y, \_) = y

third :: (a, b, c) -> c third (\_, \_, z) = z A tuple of patterns is itself a pattern, which matches any tuple of the same arity whose components all match the corresponding patterns in order.

Functions fst and snd are defined in the module Data.Tuple:

```
\begin{array}{l} \lambda > ::type \mbox{ fst } \\ \mbox{ fst } :: \ (a, \ b) \ -> \ a \\ \lambda > \ \mbox{ fst } \ (1,2) \\ 1 \\ \lambda > :type \ \mbox{ snd } \\ \mbox{ snd } :: \ (a, \ b) \ -> \ b \\ \lambda > \ \mbox{ snd } \ (1,2) \end{array}
```

```
2
```
A list of patterns is itself a pattern, which matches any list of the same length whose components all match the corresponding patterns in order.

-- three characters beginning with the letter 'a'
test :: [Char] -> Bool
test ['a',\_,\_] = True
test \_ = False

-- four characters ending with the letter 'z'
test :: [Char] -> Bool
test [\_,\_,\_,'z'] = True
test \_ = False

#### Pattern matching – Lambda expression

- An anonymous function is a function without a name.
- It is a Lambda abstraction and might look like this:
   \x -> x + 1.

(That backslash is Haskell's way of expressing a  $\lambda$  and is supposed to look like a Lambda.)

```
\lambda > :type (\x -> x+1)
(\x -> x+1) :: Num a => a -> a
\lambda > (\x -> x+1) 2
```

3

The definition

add :: Int -> Int -> Int -> Int add x y z = x+y+z

can be understood as meaning

add :: Int  $\rightarrow$  Int  $\rightarrow$  Int  $\rightarrow$  Int add =  $x \rightarrow (y \rightarrow (x \rightarrow x+y+z))$ 

which makes precise that add is a function that takes an integer x and returns a function which in turn takes another integer y and returns a function which in turn takes another integer z and returns the result x+y+z.

 $\lambda$ -expressions are useful when defining functions that returns function as results by their very nature, rather than a consequence of currying.

const :: a -> b -> a
const x \_ = x
-- emphasis const :: a -> (b -> a)
const :: a -> b -> a
const x = \\_ -> x

# Pattern matching – Lambda expression

A closure (the opposite of a combinator) is a function that makes use of free variables in its definition. It closes around some portion of its environment.

f :: Num a  $\Rightarrow$  a  $\Rightarrow$  a  $\Rightarrow$  a f x =  $y \Rightarrow x + y$ 

f returns a closure, because the variable x, which is bounded outside of the lambda abstraction is used inside its definition.

$$\begin{aligned} \lambda &> g = f 1 \\ \lambda &> g 2 \\ 3 \\ \lambda &> g 3 \\ 4 \end{aligned}$$

#### Pattern matching – Operator sections

- Functions such as + that are written between their two arguments are called section
- Any operator can be converted into a curried function by enclosing the name of the operator in parentheses, such as (+) 1 2.
- More generally, if o is an operator, then expression of the form
   (o), (x o) and (o y) are called sections whose meaning as
   functions can be formalised using λ-expressions as follows:

- (+) is the addition function  $x \rightarrow (y \rightarrow x+y)$ .
- (1 +) is the successor function  $y \rightarrow 1+y$ .
- (1 /) is the reciprocation function  $y \rightarrow 1/y$ .
- (\* 2) is the doubling function  $x \rightarrow x*2$ .
- (/ 2) is the halving function  $x \rightarrow x/2$ .

- A where clause is used to divide the more complex logic or calculation into smaller parts, which makes the logic or calculation easy to understand and handle
- A where clause is bound to a surrounding syntactic construct, like the pattern matching line of a function definition.
- A where clause is a syntactic construct

# Pattern matching – Bindings

```
bmiTell :: (RealFloat a) => a -> a -> String
bmiTell weight height
  | weight / height ^ 2 <= 18.5 = "Underweight"
  | weight / height ^ 2 < 25.0 = "Healthy weight"
  | weight / height ^ 2 < 30.0 = "Overweight"
  | otherwise = "Obese"
```

```
bmiTell :: (RealFloat a) => a -> a -> String
bmiTell weight height
| bmi <= 18.5 = "Underweight"
| bmi < 25.0 = "Healthy weight"
| bmi < 30.0 = "Overweight"</pre>
```

| otherwise = "Obese"

#### where

bmi = weight / height ^ 2

```
bmiTell :: (RealFloat a) => a -> a -> String
bmiTell weight height
  | bmi <= underweight = "Underweight"
  bmi < healthy = "Healthy weight"</pre>
  bmi < overweight = "Overweight"</pre>
                = "Obese"
  | otherwise
  where
    bmi = weight / height ^ 2
   underweight = 18.5
   healthy = 25
```

overweight = 30

- A let binding binds variables anywhere and is an expression itself, but its scope is tied to where the let expression appears.
- if a let binding is defined within a guard, its scope is local and it will not be available for another guard.
- A let binding can take global scope overall pattern-matching clauses of a function definition if it is defined at that level.

```
cylinder :: (RealFloat a) => a -> a -> a
cylinder r h =
  let sideArea = 2 * pi * r * h
     topArea = pi * r ^2
  in sideArea + 2*topArea
```

#### Pattern matching – Bindings

```
\lambda > let zoot x y z = x*y + z
\lambda > :type zoot
zoot :: Num a => a -> a -> a -> a
\lambda > zoot 3 9 2
29
\lambda > let boot x y z = x*y + z in boot 3 9 2
29
\lambda > :type boot
<interactive>: error:
  o Variable not in scope: boot
```

#### Pattern matching – Bindings

 $\lambda$  > let a = 1; b = 2 in a + b 3  $\lambda$  > let a = 1; b = a + 2 in a + b 4  $\lambda$  > let a = 1; a = 2 in a <interactive>:: error: Conflicting definitions for 'a'  $\lambda$  > let a = 1; b = 2+a; c = 3+a+b in (a, b, c) (1,3,7)

$$\lambda$$
 > let a = 1 in let a = 2; b = 3+a in b  
5

 $\lambda$  > let a = 1 in let a = a+2 in let b = 3+a in b ^CInterrupted.

$$\lambda$$
 > let f x y = x+y+1 in f 3 5  
9

 $\lambda>$  let f x y = x+y; g x = f x (x+1) in g 5 11

#### Pattern matching – Bindings

```
dist :: Floating a => (a, a) -> (a, a) -> a
dist (x1,y1) (x2,y2) =
    let xdist = x2 - x1
        ydist = y2 - y1
        sqr z = z*z
        in sqrt ((sqr xdist) + (sqr ydist))
```

dist :: Floating a => (a, a) -> (a, a) -> a
dist (x1,y1) (x2,y2) = sqrt((sqr xdist) + (sqr ydist))
where

We can pattern match with let bindings. E.g., we can dismantle a tuple into components and bind the components to names.

 $\lambda$  > f x y z = let (sx,sy,sz) = (x\*x,y\*y,z\*z) in (sx,sy,sz)  $\lambda >$  f 1 2 3 (1, 4, 9) $\lambda$  > g x y = let (sx,\_) = (x\*x,y\*y) in sx  $\lambda > g 2 3$ 4  $\lambda > h x = let ((sx, cx), qx) = ((x + x, x + x + x), x + x + x + x) in (sx, cx, qx)$  $\lambda > h 2$ (4, 8, 16)

# Pattern matching – Bindings

```
let bindings are expressions.
\lambda > 1 + \text{let } x = 2 \text{ in } x + x
5
\lambda > (let x = 2 in x*x) + 1
5
\lambda > (let (x,y,z) = (1,2,3) in x+y+z) * 100
600
\lambda > (let x = 2 in (+ x)) 3
5
\lambda > let x=3 in x*x + let x=4 in x*x
25
```

# **Some functions**

Functional programming concepts

First steps

Types and classes

Defining functions

Some functions

The double factorial (or semifactorial of a number n, denoted by n!!, is the product of all the integers from 1 up to n that have the same parity (odd or even) as n

```
The double factorial (or semifactorial of a number n, denoted by
n!!, is the product of all the integers from 1 up to n that have the
same parity (odd or even) as n
dblFact1 :: Integral a => a -> a
dblFact1 n = go n
  where
    p m = (even n \&\& even m) || (odd n \&\& odd m)
    go 0 = 1
    go m
       | p m = m * go (m-1)
```

| otherwise = go (m-1)

```
The double factorial (or semifactorial of a number n, denoted by n!!, is the product of all the integers from 1 up to n that have the same parity (odd or even) as n
```

```
dblFact2 :: Integral a => a -> a
dblFact2 n = go n
 where
   nParity2 = n `mod` 2
   p m = m `mod` 2 == nParity2
   go 0 = 1
   go m
      | pm = m * go (m-1)
      | otherwise = go (m-1)
```

The double factorial (or semifactorial of a number n, denoted by n!!, is the product of all the integers from 1 up to n that have the same parity (odd or even) as n

dblFact3 :: (Eq a, Num a) => a -> a dblFact3 0 = 1 dblFact3 1 = 1 dblFact3 n = n \* dblFact3 (n-2) The double factorial (or semifactorial of a number n, denoted by n!!, is the product of all the integers from 1 up to n that have the same parity (odd or even) as n

```
dblFact4 :: (Num a, Enum a) => a -> a
dblFact4 n = product [n,n-2..1]
```

# **Collatz conjecture**

The Collatz conjecture is one of the most famous unsolved problems in mathematics. It concerns sequences of integers in which each term is obtained from the previous term as follows:

$$u_n = \begin{cases} u_{n-1}/2 & \text{if } u_{n-1} \text{ is even} \\ 3u_{n-1} + 1 & \text{if } u_{n-1} \text{ is odd} \end{cases}$$

For instance, starting with n = 19, one gets the sequence 19, 58, 29, 88, 44, 22, 11, 34, 17, 52, 26, 13, 40, 20, 10, 5, 16, 8, 4, 2, 1.

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For instance, starting with n = 19, one gets the sequence 19, 58, 29, 88, 44, 22, 11, 34, 17, 52, 26, 13, 40, 20, 10, 5, 16, 8, 4, 2, 1.

```
collatz2 :: Integral a => a -> String
collatz2 1 = "win"
collatz2 n
  | even n = collatz2 (n `div` 2)
  | otherwise = collatz2 (3*n + 1)
```

$$A(0, n) = n + 1$$
  

$$A(m + 1, 0) = A(m, 1)$$
  

$$A(m + 1, n + 1) = A(m, A(m + 1, n))$$

```
aP :: (Num a, Eq a, Num b, Eq b) => a -> b -> b

aP 0 n = n+1

aP m 0 = aP (m-1) 1

aP m n = aP (m-1) (aP m (n-1))
```

A prime number (or a prime) is a natural number greater than 1 that is not a product of two smaller natural numbers.

```
-- odd number predicate

isOdd :: (Eq a, Num a) => a -> Bool

isOdd 0 = False

isOdd 1 = True

isOdd n = isEven (n-1)
```

```
-- even number predicate

isEven :: (Eq a, Num a) => a -> Bool

isEven 0 = True

isEven 1 = False

isEven n = isOdd (n-1)
```

fact1 :: (Eq a, Num a) => a -> a fact1 n = if n == 0 then 1 else n \* fact1 (n-1) fact2 :: (Eq a, Num a) => a -> a

fact2 n

```
| n == 0 = 1
| otherwise = n * fact2 (n-1)
```

#### Factorial

```
fact3 :: (Ord a, Num a) \Rightarrow a \Rightarrow a
fact3 = go 1
  where
    go m n
       | m > n = 1
       | otherwise = m * go (m+1) n
fact4 :: (Eq t, Num t) \Rightarrow t \rightarrow t
fact4 n = go 1 n
  where
    go acc 0 = acc
    go acc m = go (acc*m) (m-1)
```

```
fact5 :: (Enum a, Num a) => a -> a
fact5 n = product [1..n]
```

# Pascal triangle

1 1 1 1 2 1 3 3 1 1 4 6 1 4 1 5 1 5 10 10 1 1 6 15 20 15 6 1 1 7 21 35 35 21 7 1 pT :: (Num a, Ord a, Num b) => a -> a -> b pTrc | r == 1 && c == 1 = 1| c < 1 || c > r = 0| otherwise = pT (r-1) (c-1) + pT (r-1) c